



Surface Triangulation and Voronoi Regions

CS334
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[Slides with help from Michael Kazhdan @ JHU,
Ioannis Stamos @ CUNY, and
Prof. Shmuel Wimer and Andy Mirzaian]



Motivation

- Time of flight
- Structured light
- Stereo images
- Shape from shading
- Etc.

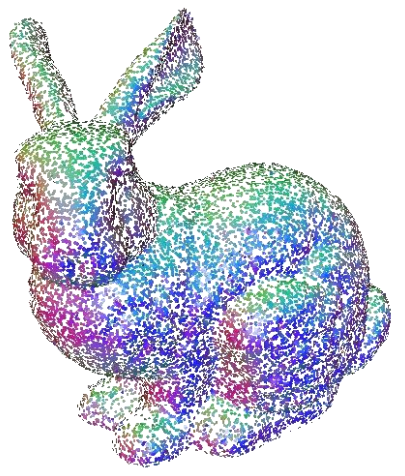
■ <http://graphics.stanford.edu/projects/mich/>



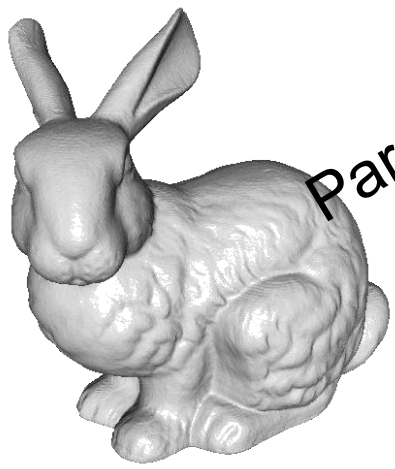


Motivation

Surface reconstruction



Geometry processing

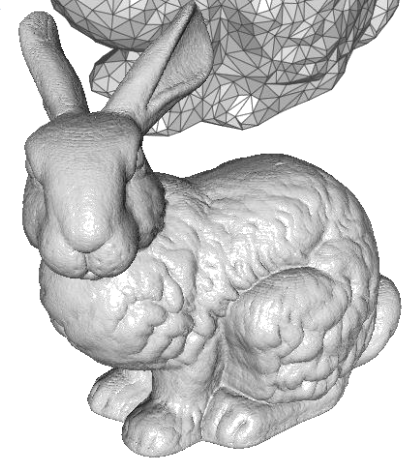
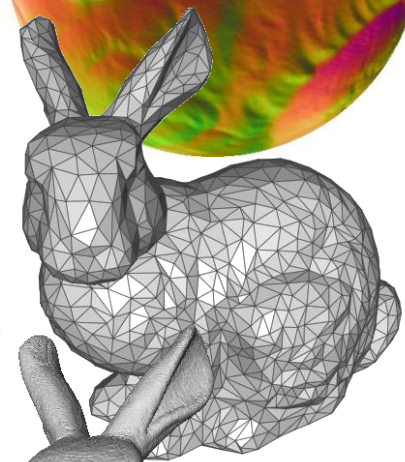
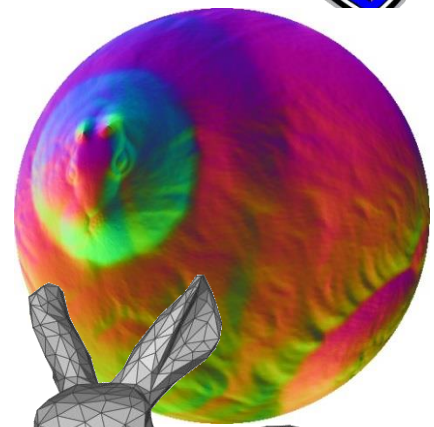


Parameterization

Decimation

Filtering

etc.

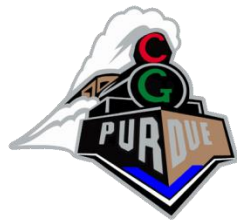


Marching Squares (2D)

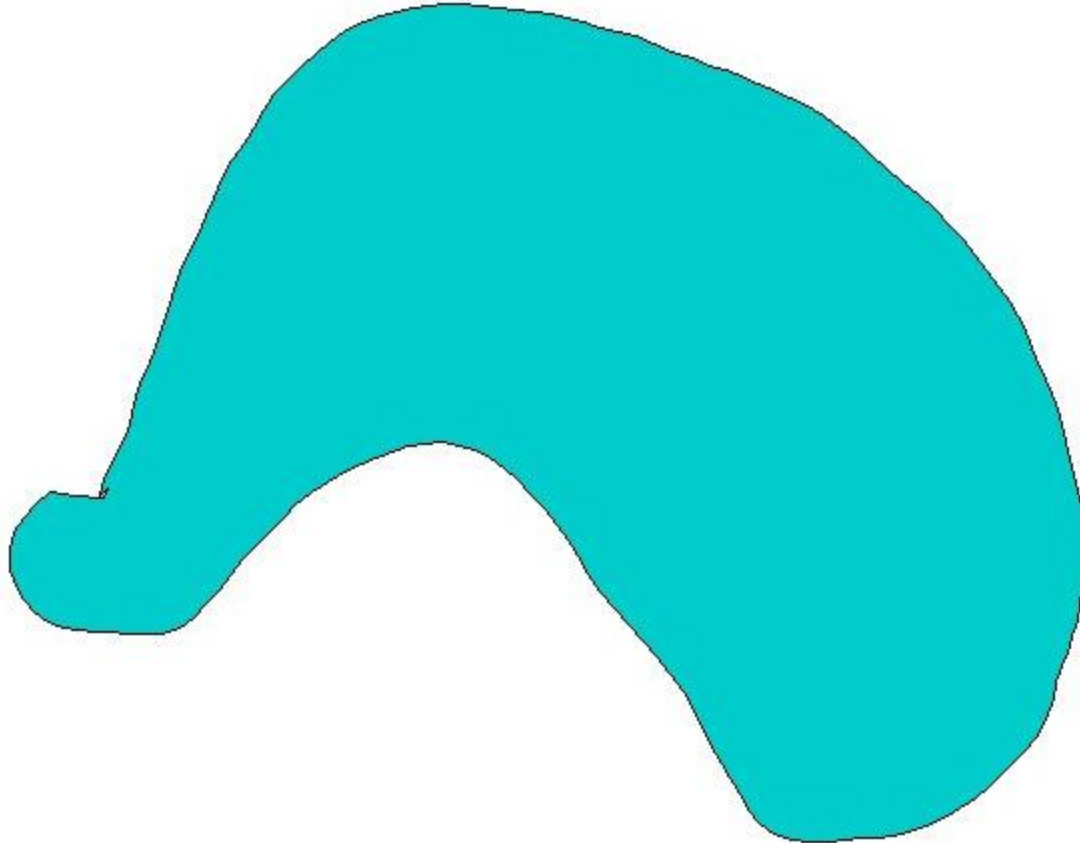
Marching Cubes (3D)



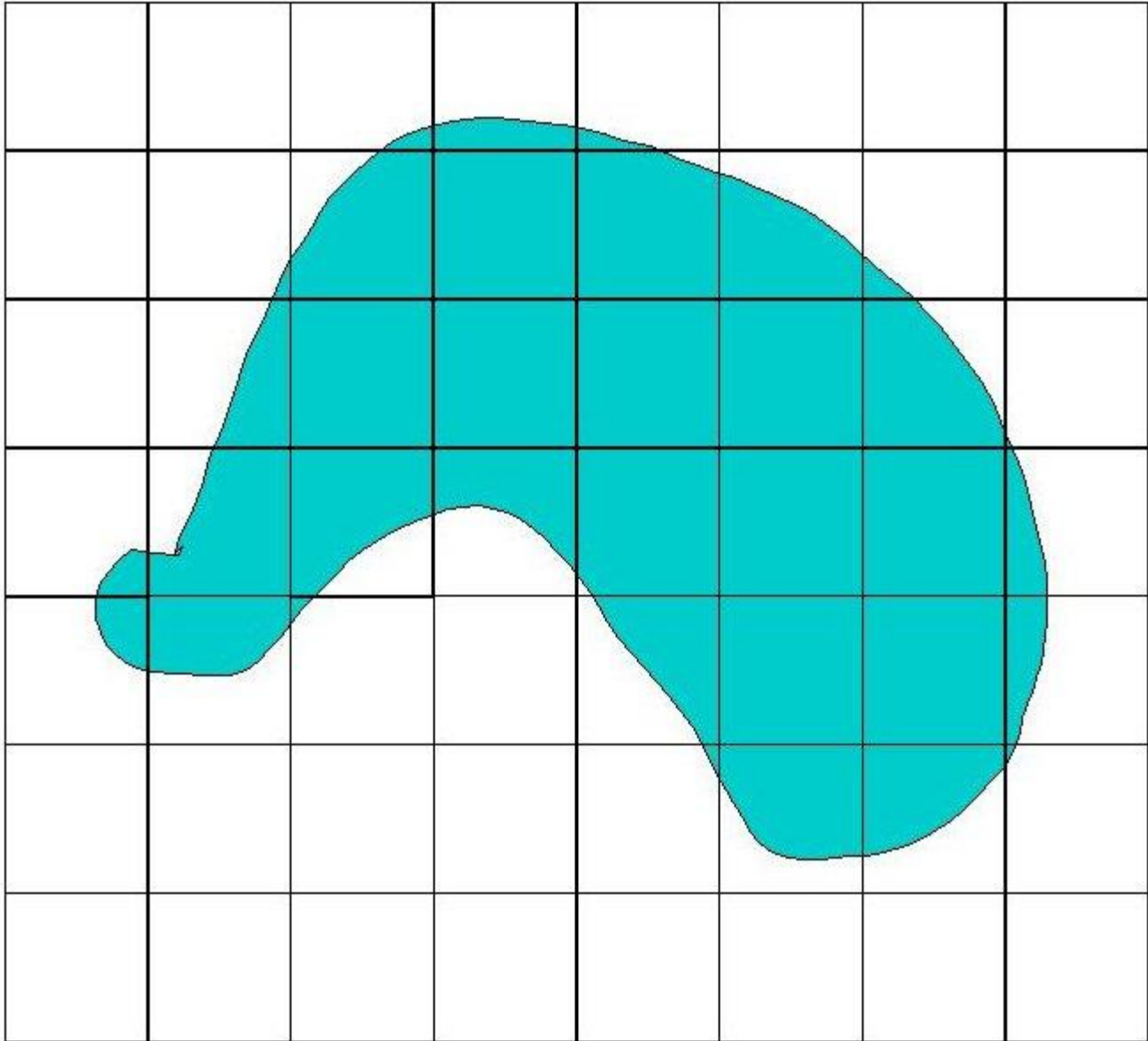
- The premise of the algorithm is to divide the input volume into a discrete set of squares/cubes. By assuming linear reconstruction filtering, each square/cube, which contains a piece of a given isosurface, can be easily represented with lines/triangles.

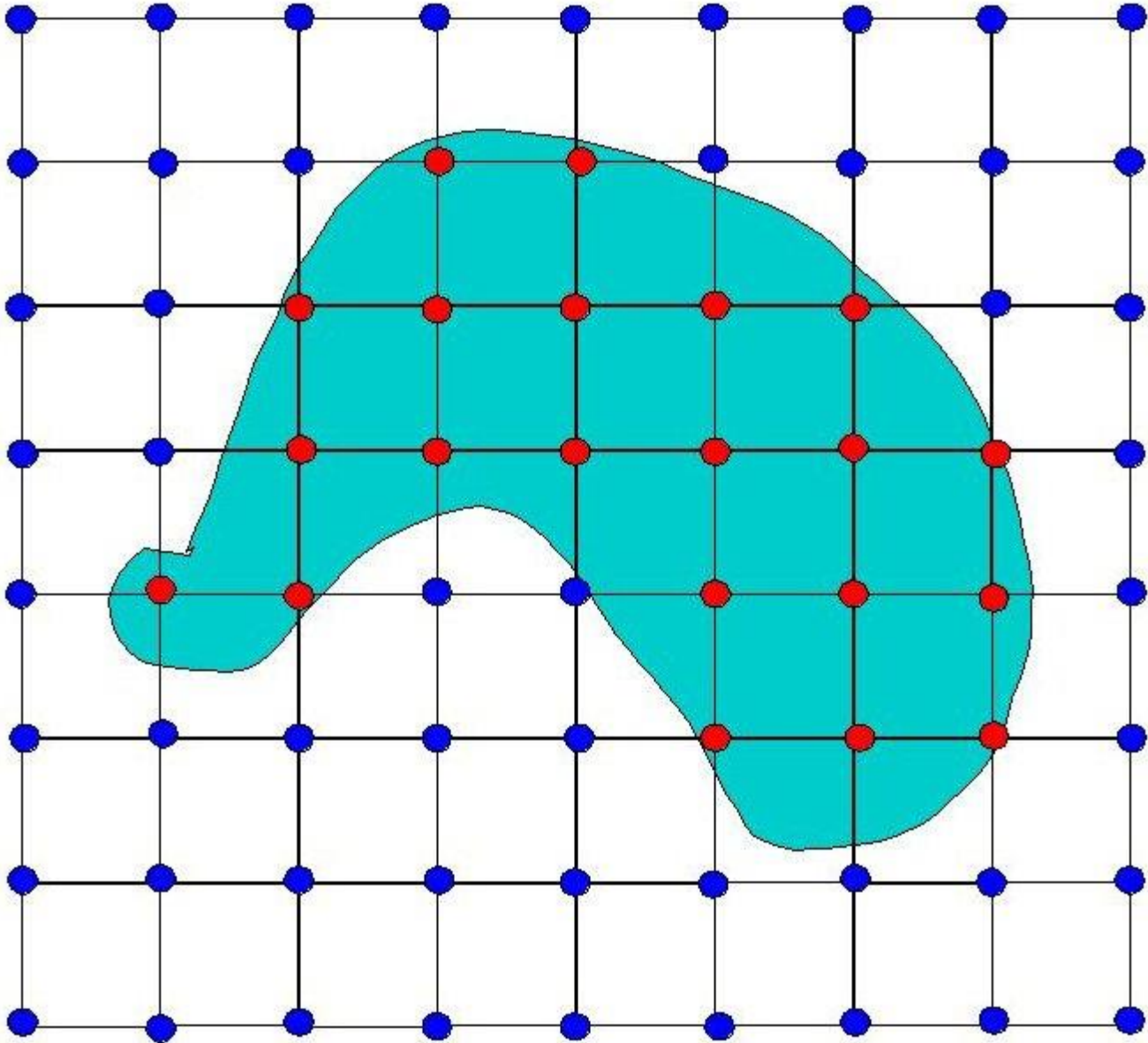


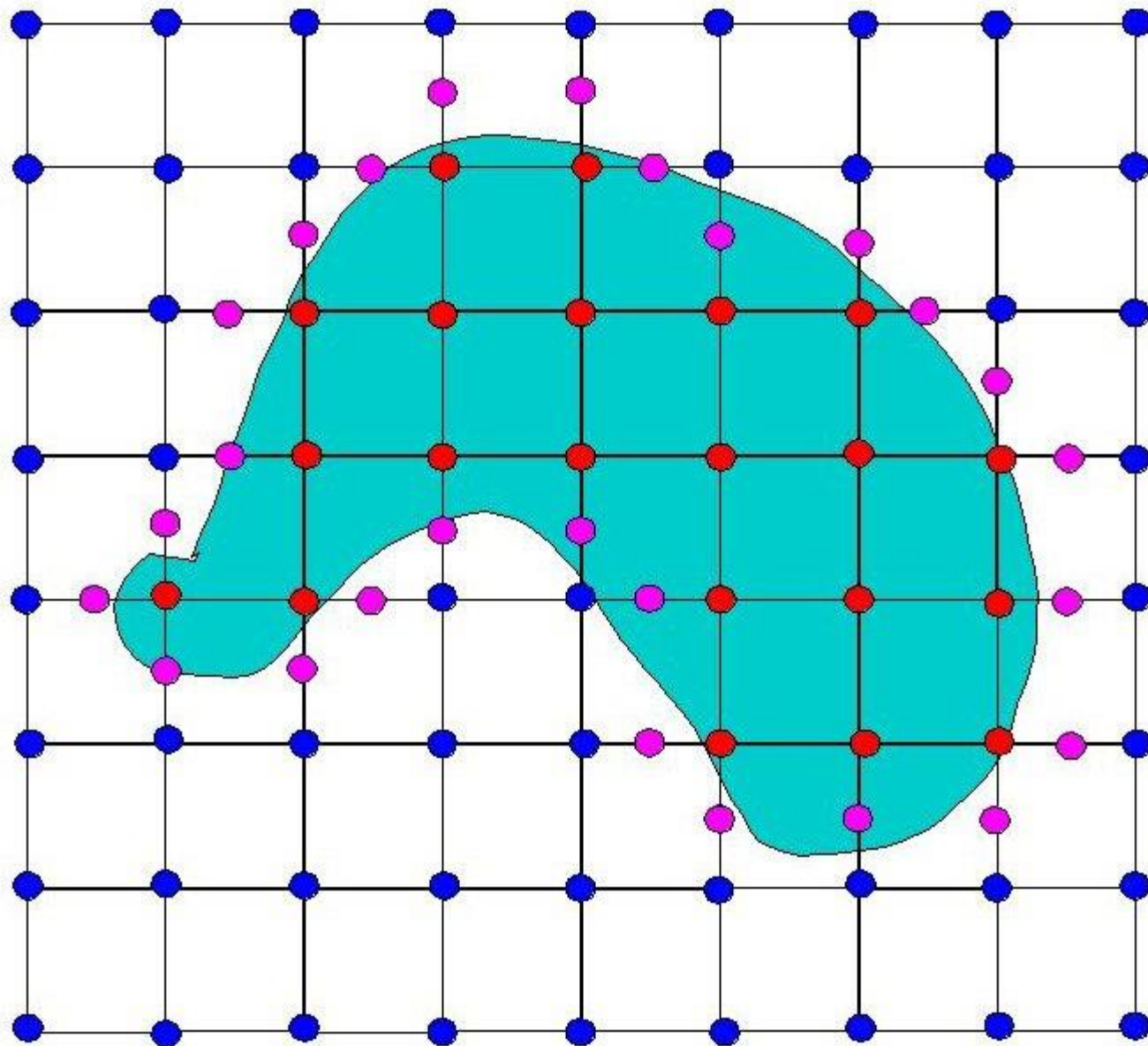
Marching Squares

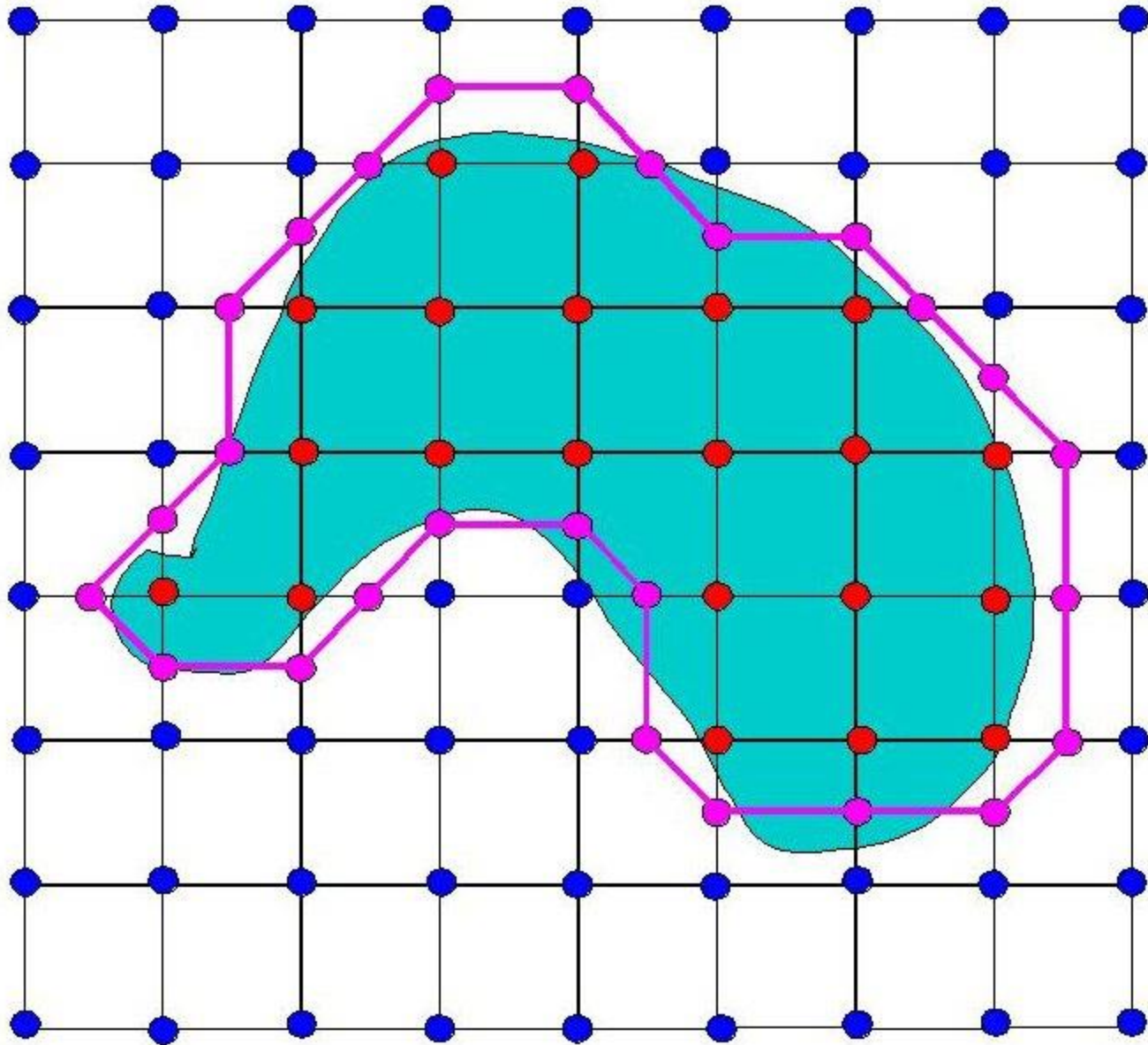


[https://www.cs.carleton.edu/cs_comps/0405/shape/marching_cubes.html]





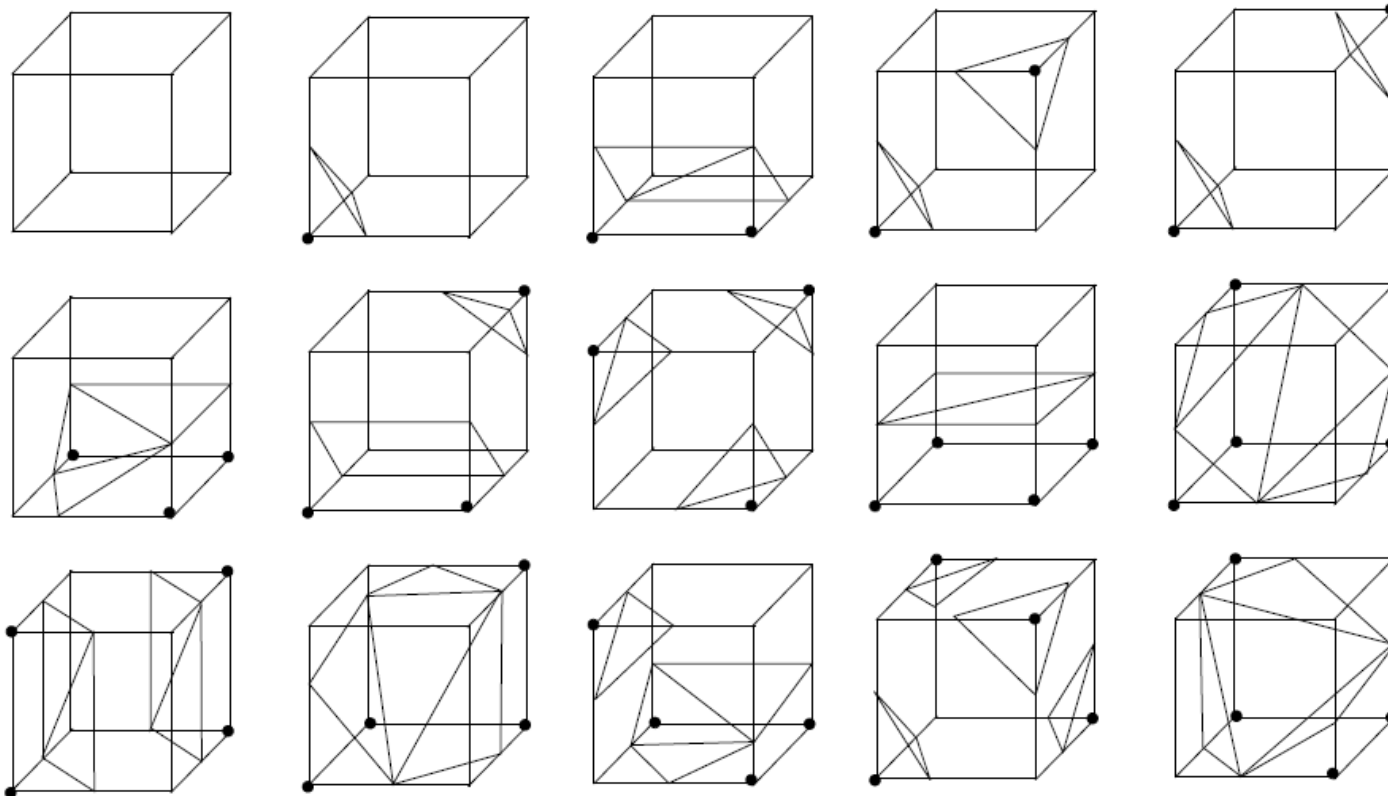






Marching Cubes

If the function is sampled on a regular voxel grid, we can independently triangulate each voxel.

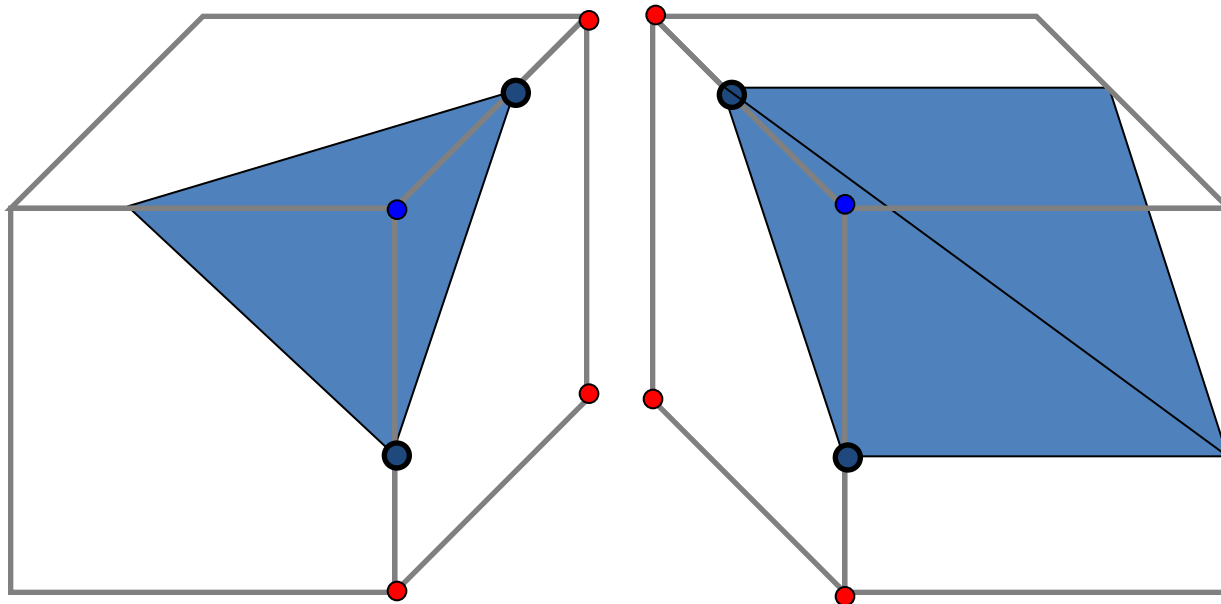




Marching Cubes

Iso-vertices on an edge are only determined by the values on the corner of the edge:

⇒ Iso-vertices are consistent across voxels.

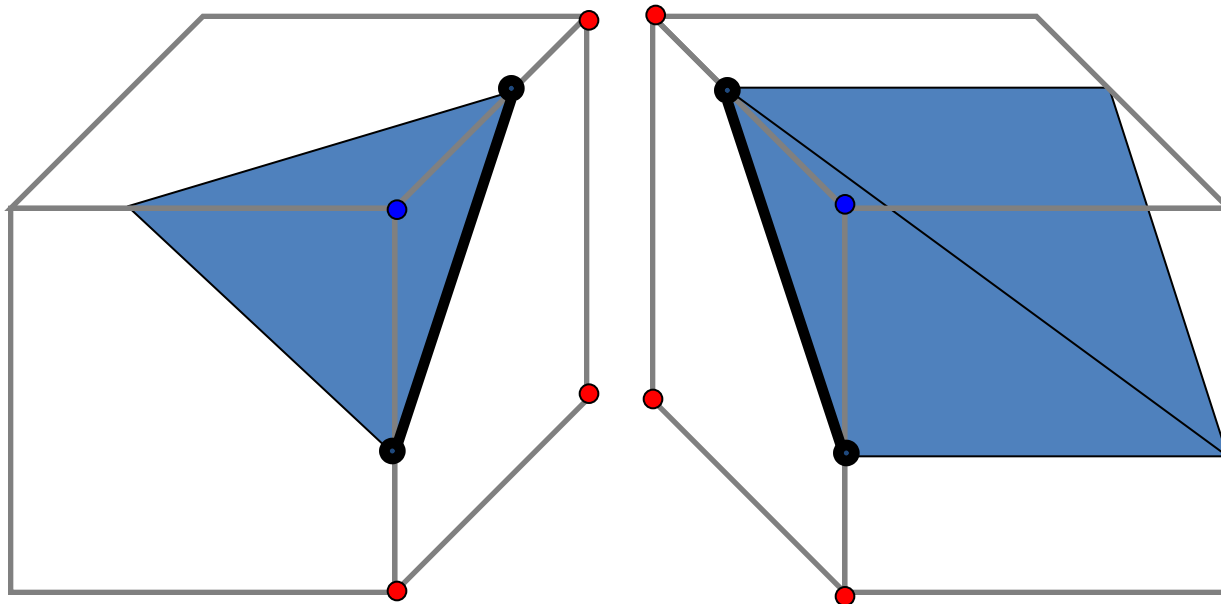


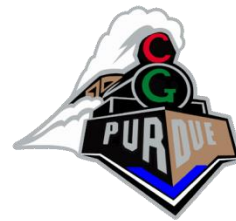


Marching Cubes

Iso-edges on a face are only determined by the values on the face:

⇒ Each iso-edge is shared by two triangles so the mesh is water-tight.

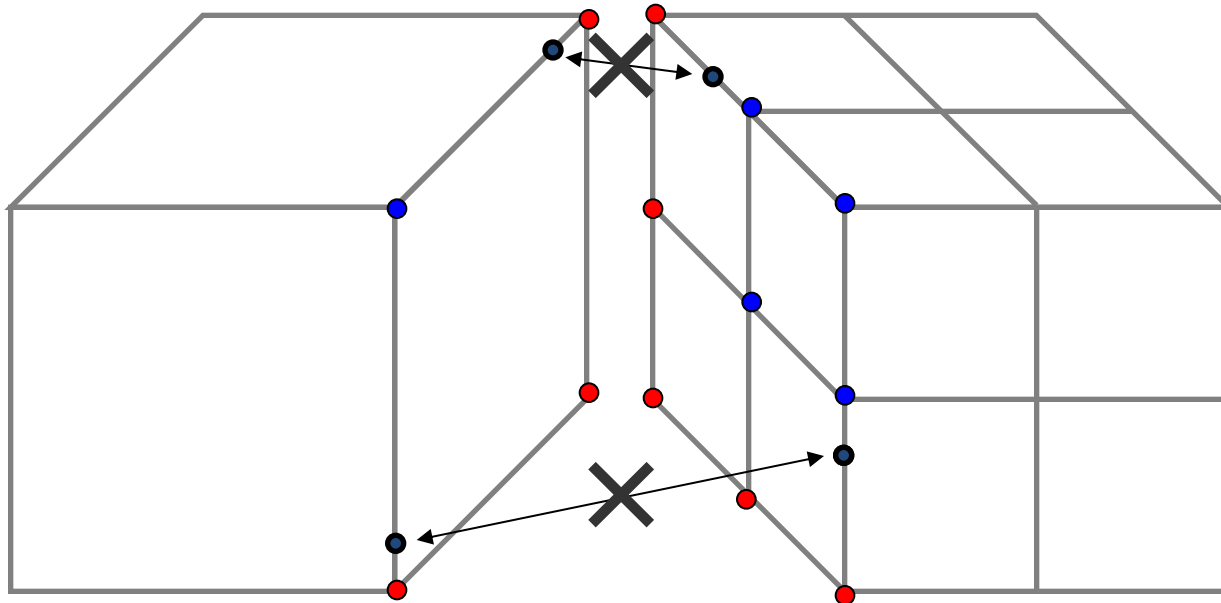


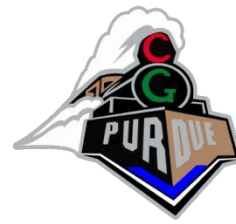


Challenges

Extracting a surface by independently triangulating the leaf octants, depth-disparities can cause:

- Inconsistent extrapolation to edges
⇒ Inconsistent iso-vertex positions

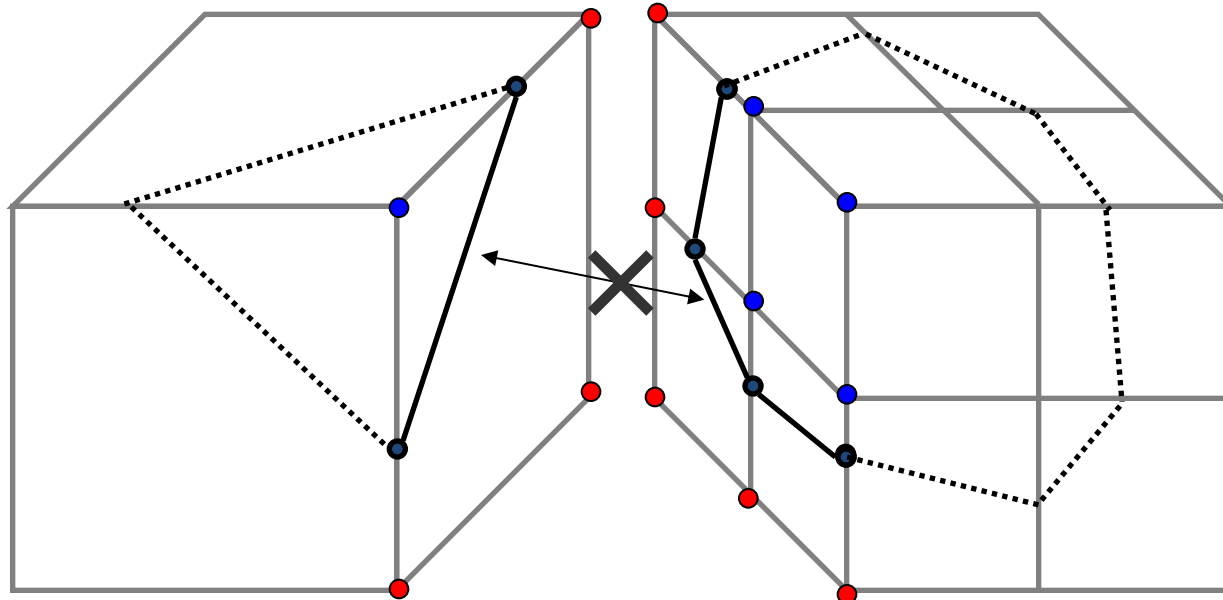




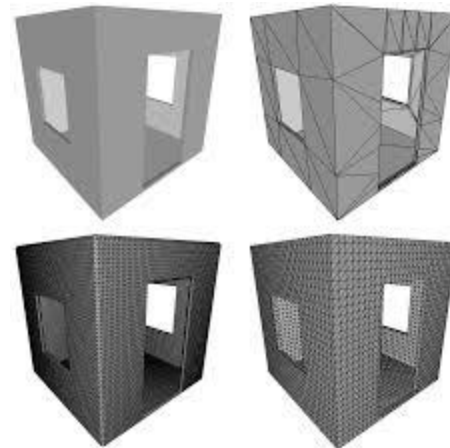
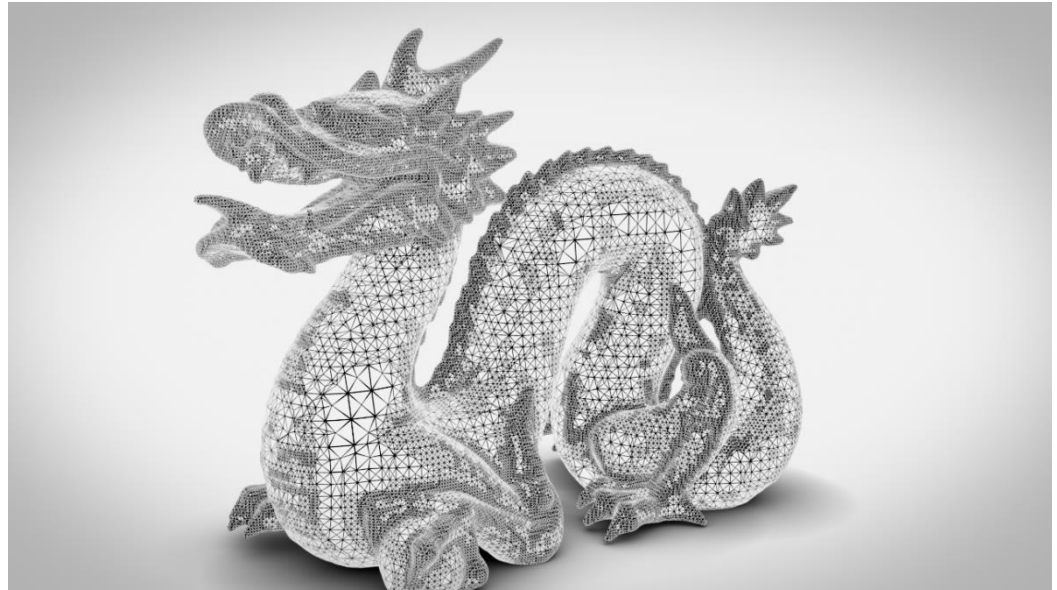
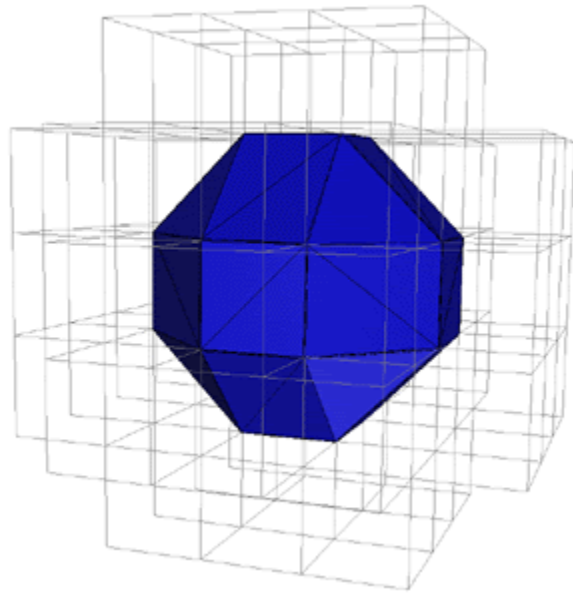
Challenges

Extracting a surface by independently triangulating the leaf octants, depth-disparities can cause:

- Inconsistent extrapolation to faces
⇒ Inconsistent iso-edges



Marching Cubes

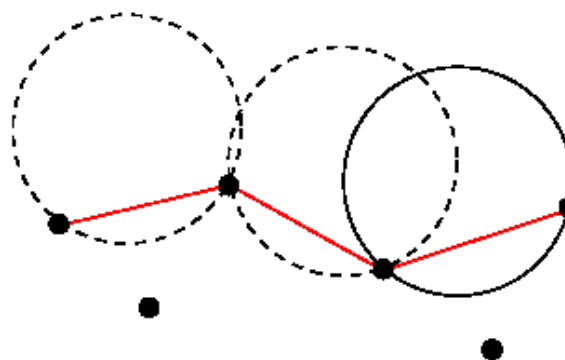




Ball-pivoting



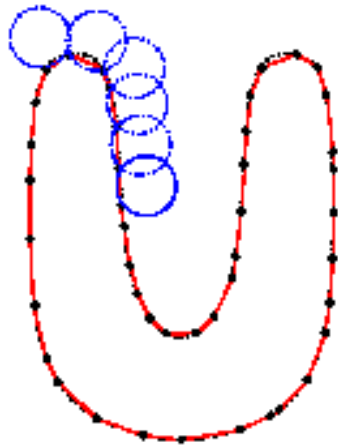
Bernardini et al., IBM



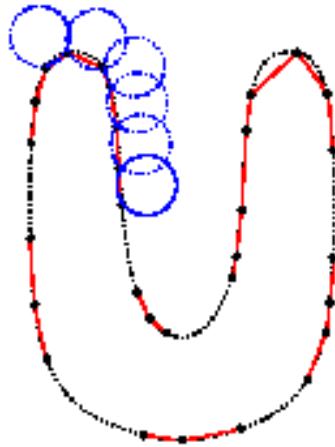
Fixed-radius ball “rolling”
over points selects subset of
alpha-shape.



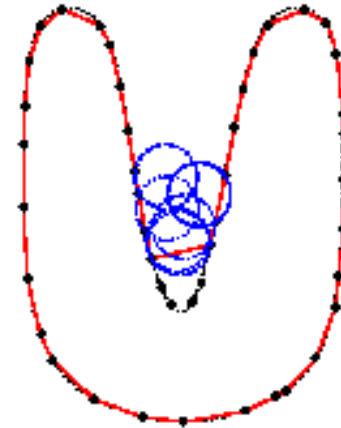
Pivoting in 2D



(a)



(b)

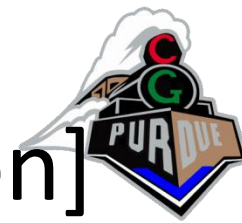


(c)

(a) Circle of radius ρ pivots from point to point, connecting them with edges.

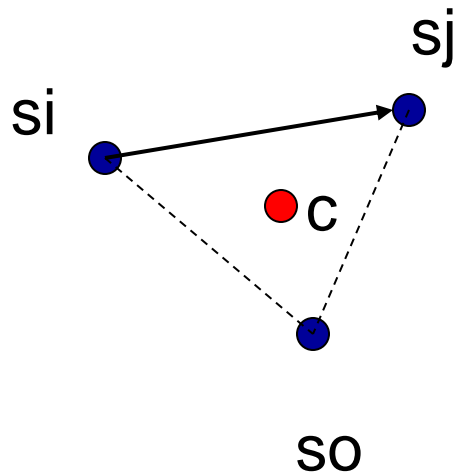
(b) When sampling density is low, some of the edges will not be created, leaving holes.

■ (c) When the curvature of the manifold is larger than $1/\rho$, some of the points will not be reached by the pivoting ball, and features will be missed.



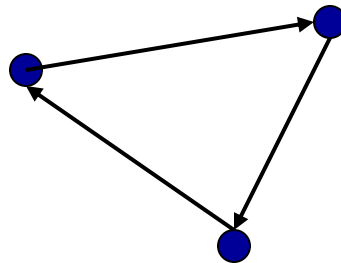
The algorithm [Edge representation]

- Edge (s_i, s_j)
 - Opposite point s_o , center of empty ball c
 - Edge: “Active”, “Boundary”, or “Frozen”





Pivoting example



Initial seed triangle:

Empty ball of radius ρ passes through the three points

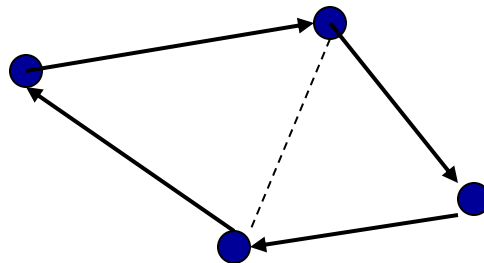
Active edge



● Point on front



Pivoting example



Ball pivoting around active edge

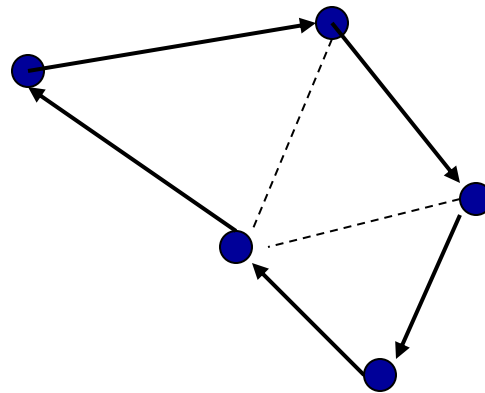
Active edge



● Point on front



Pivoting example

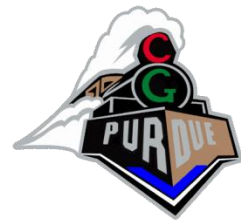


Ball pivoting around active edge

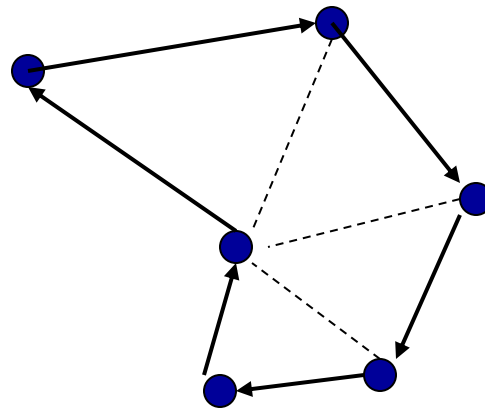
Active edge



● Point on front



Pivoting example



Ball pivoting around active edge

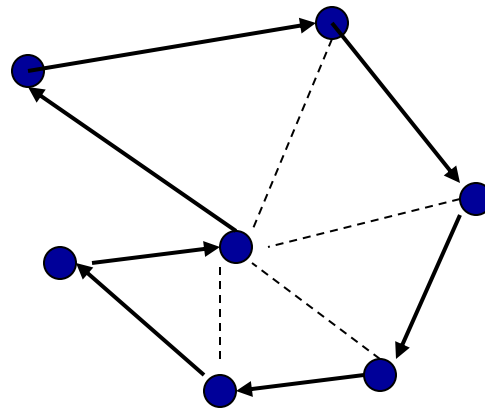
Active edge



● Point on front



Pivoting example



Ball pivoting around active edge

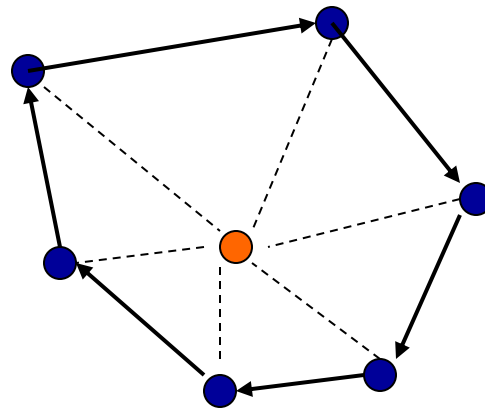
Active edge



● Point on front



Pivoting example



Ball pivoting around active edge

Active edge



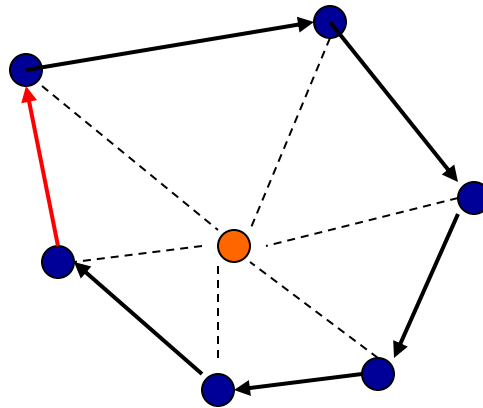
● Point on front

● Internal point



Pivoting example

Boundary edge



Ball pivoting around active edge
No pivot found

Active edge

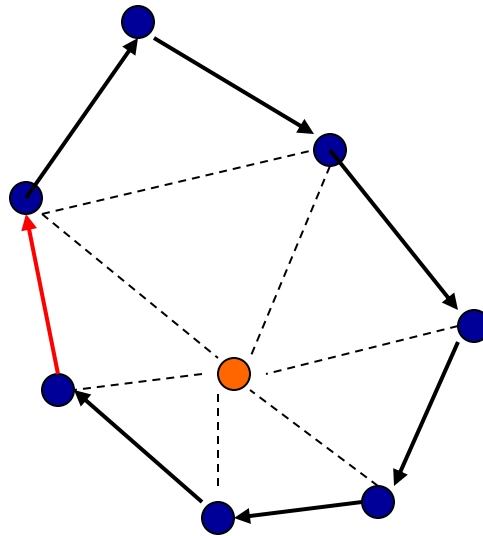


- Point on front
- Internal point



Pivoting example

Boundary edge

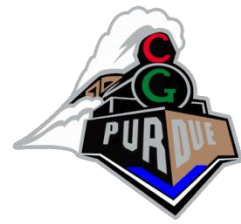


Ball pivoting around active edge

Active edge

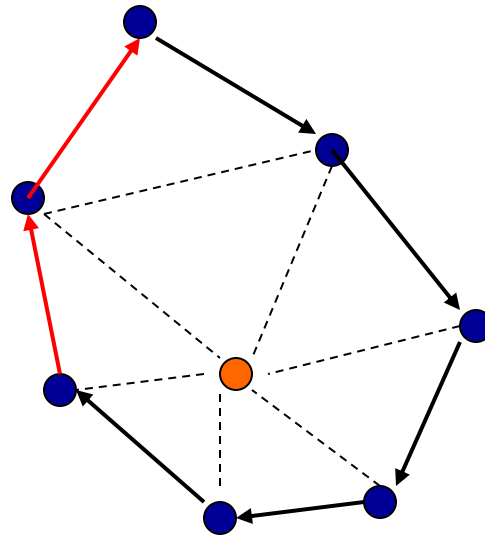


- Point on front
- Internal point



Pivoting example

Boundary edge



Ball pivoting around active edge
No pivot found

Active edge

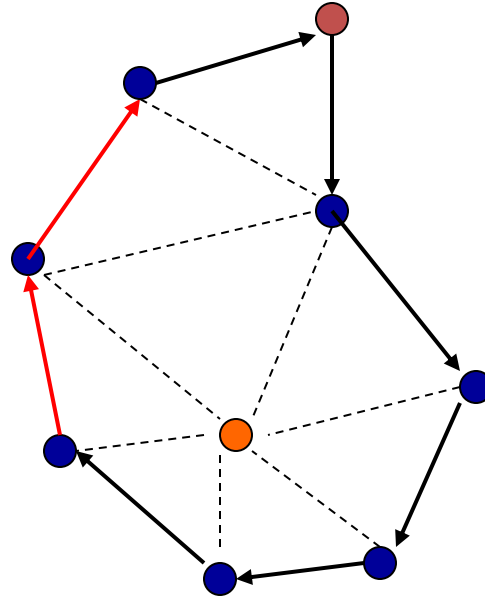


- Point on front
- Internal point



Pivoting example

Boundary edge



Ball pivoting around active edge

Active edge



- Point on front
- Internal point

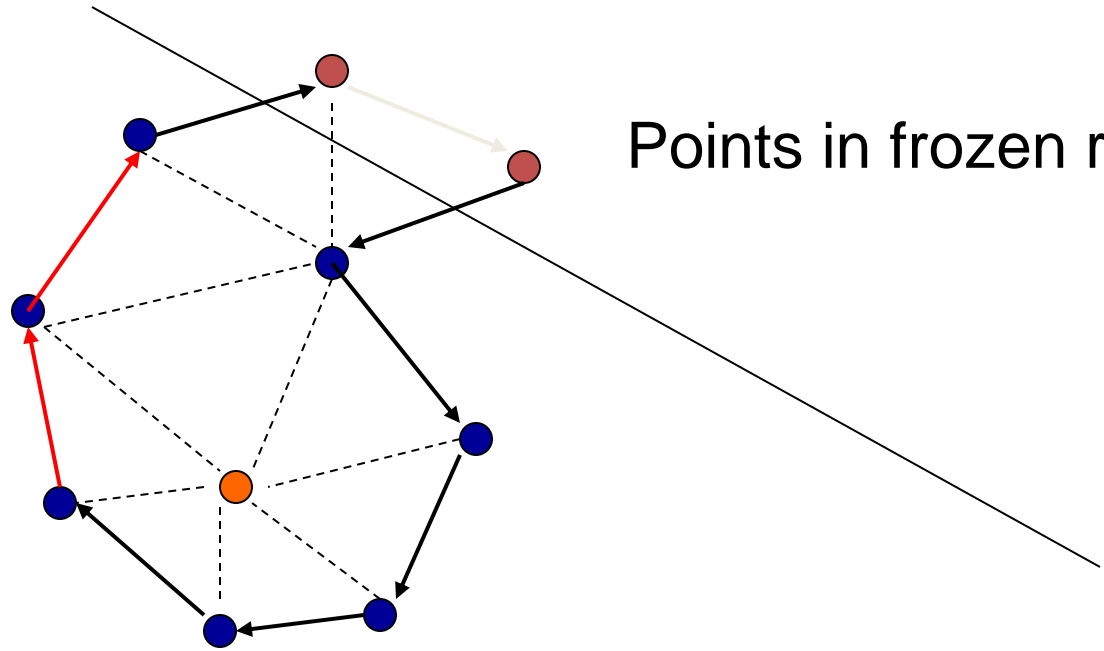


Pivoting example

Boundary edge



Frozen edge



Points in frozen region

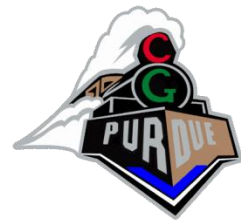
Ball pivoting around active edge

Active edge



● Point on front

● Internal point

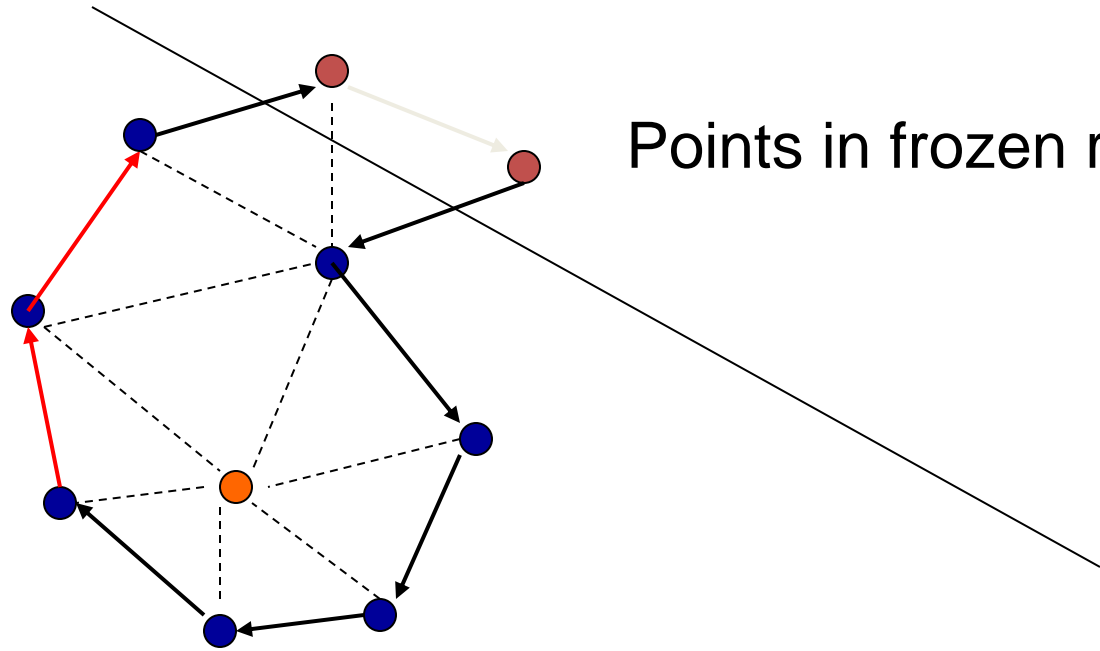


Pivoting example

Boundary edge



Frozen edge



Points in frozen region

Ball pivoting around active edge

Active edge



● Point on front

● Internal point

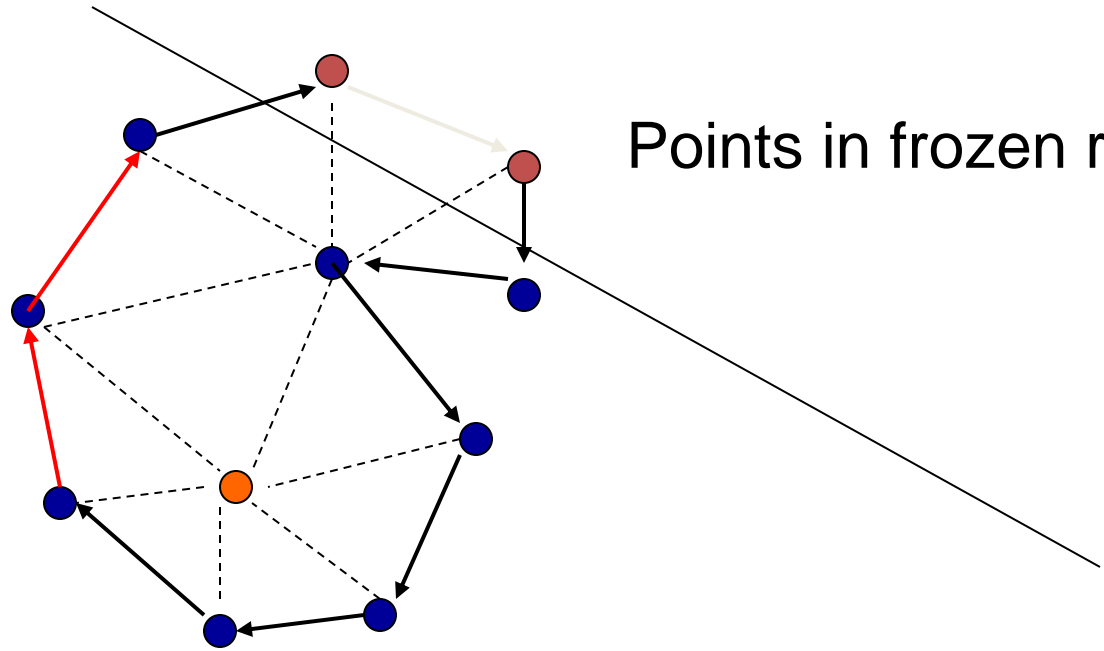


Pivoting example

Boundary edge



Frozen edge



Points in frozen region

Active edge



● Point on front

● Internal point

Ball pivoting around active edge

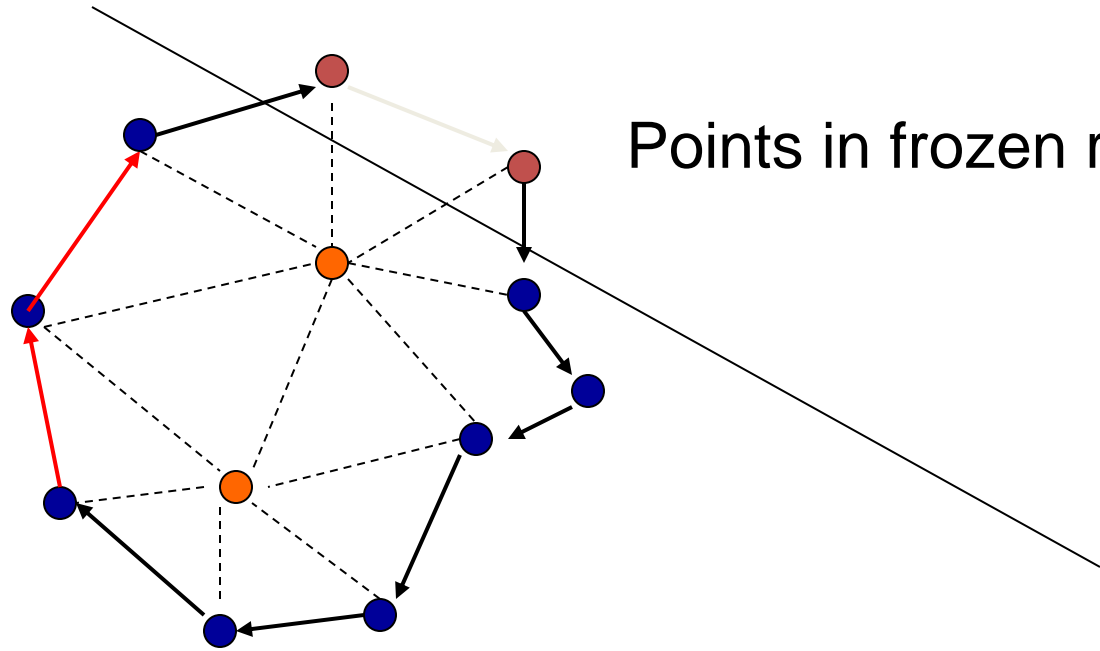


Pivoting example

Boundary edge



Frozen edge



Points in frozen region

Active edge



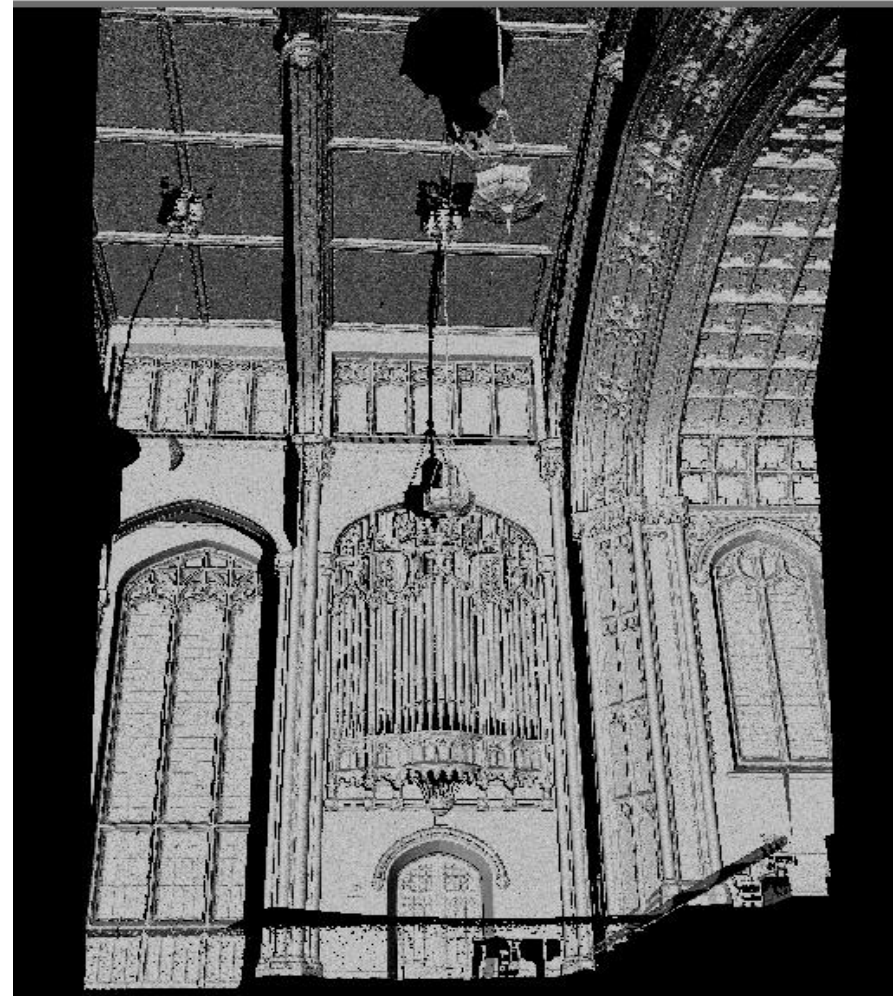
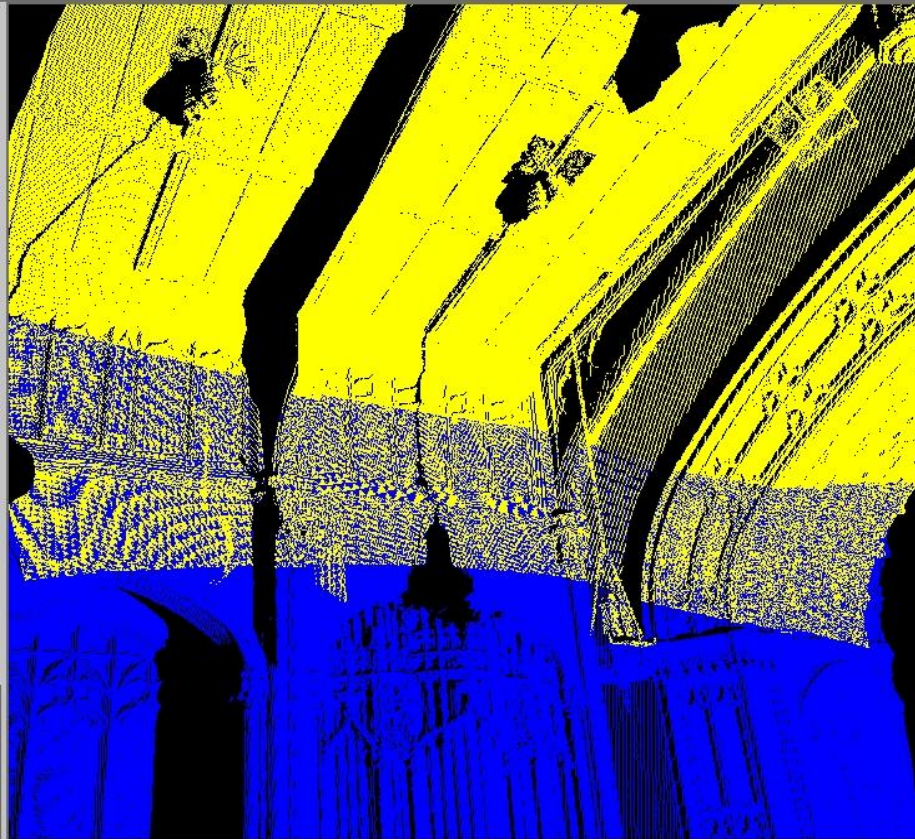
● Point on front

● Internal point

Ball pivoting around active edge



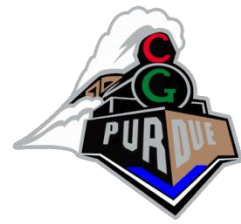
Ball Pivoting Algorithm





Ball Pivoting Algorithm

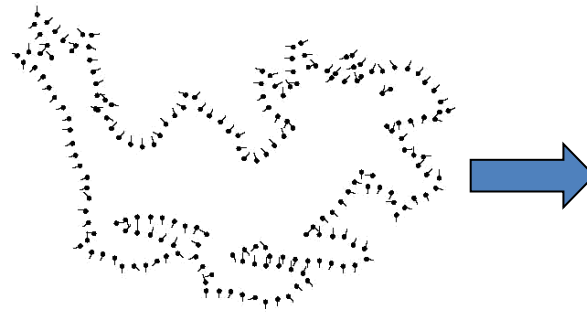




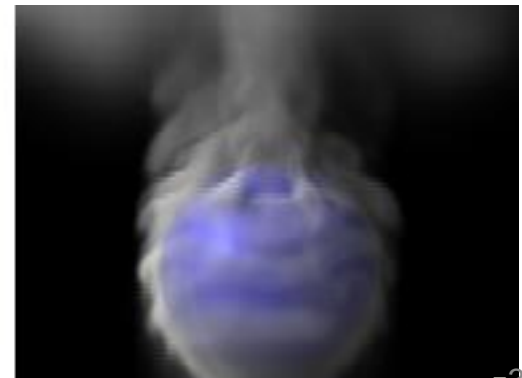
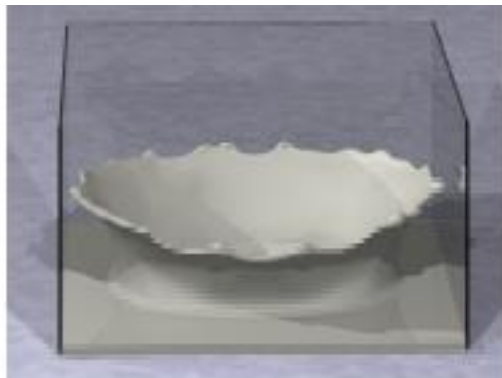
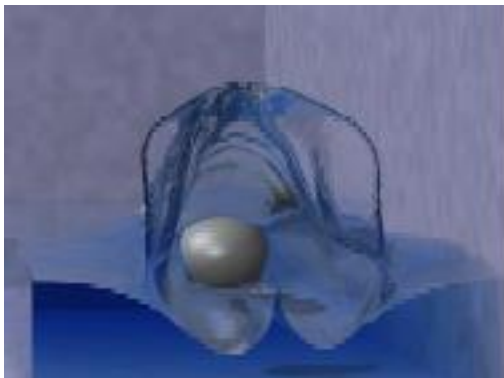
Implicit Representation

Another option is representing a 3D model by an implicit function for:

- Reconstruction
- Fluid Dynamics
- 3D Texturing



Kazhdan 2005

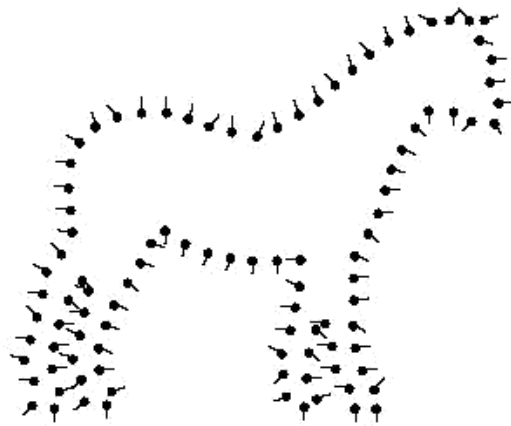




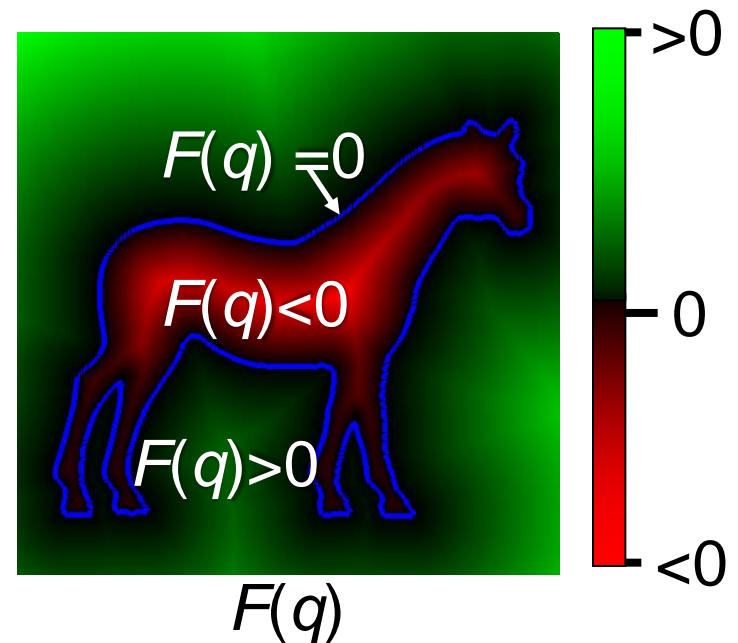
Implicit Function Fitting

Given point samples:

- Define a function with value zero at the points.
- Extract the zero isosurface.



Sample points



Triangulation Complexity (in general)



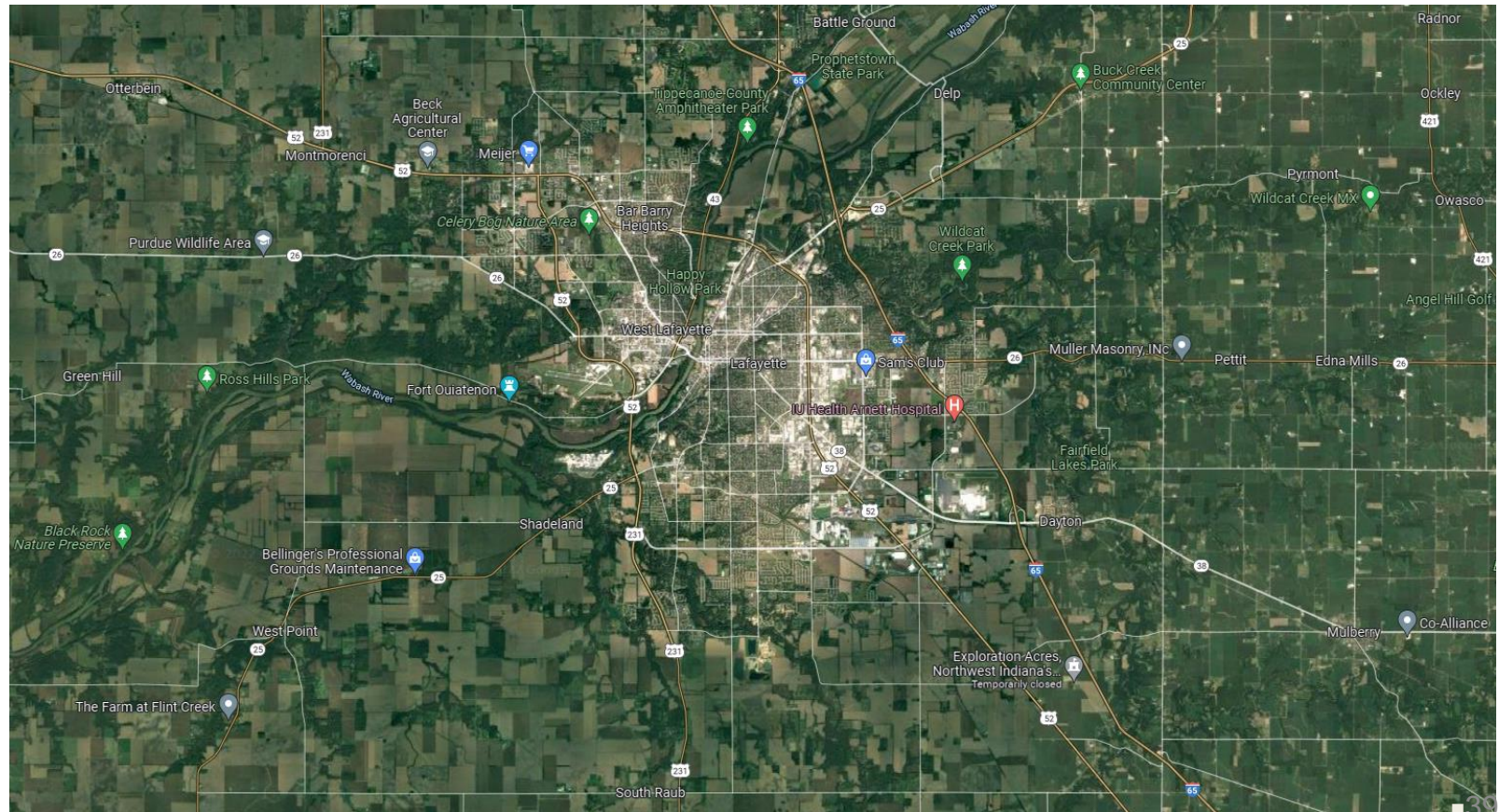
- Theorem: (Gary et. al. 1978) A simple n -vertex polygon can be triangulated in $O(n \log n)$ time and $O(n)$ storage
- The problem has been studied extensively between 1978 and 1991, when in 1991 Chazelle presented an $O(n)$ time complexity algorithm.



Delaunay Triangulation

- Another very popular algorithm...
- But first, Voronoi Diagrams...
- Relevant Conversation:
 - Captain Kirk: “Spock! Which tricorder tower (i.e., cell phone) should I be using?”
 - Commander Spock: “Logically, the closest one, Jim.”
- How do you do that?

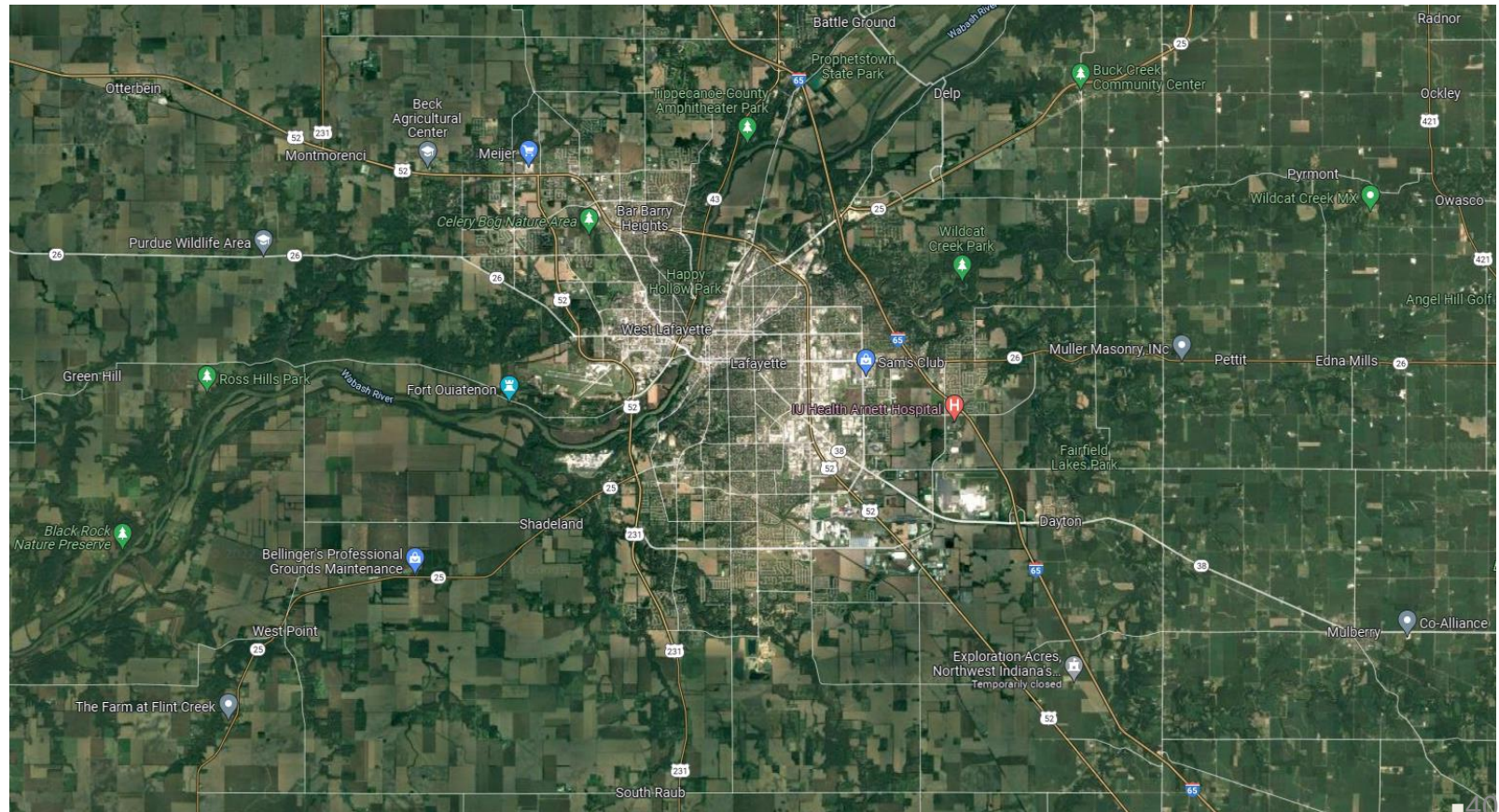
Where to place cell phone towers? or Which cell phone tower should I use?





Cell phone towers

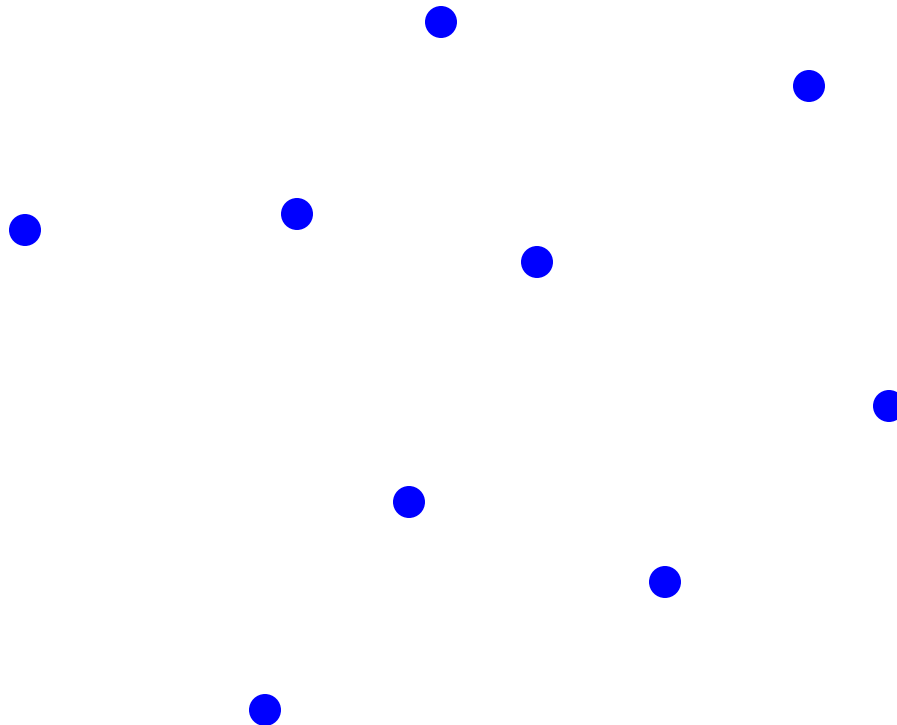
$P = \{ p_1, p_2, \dots, p_n \}$ a set of n points in the plane.





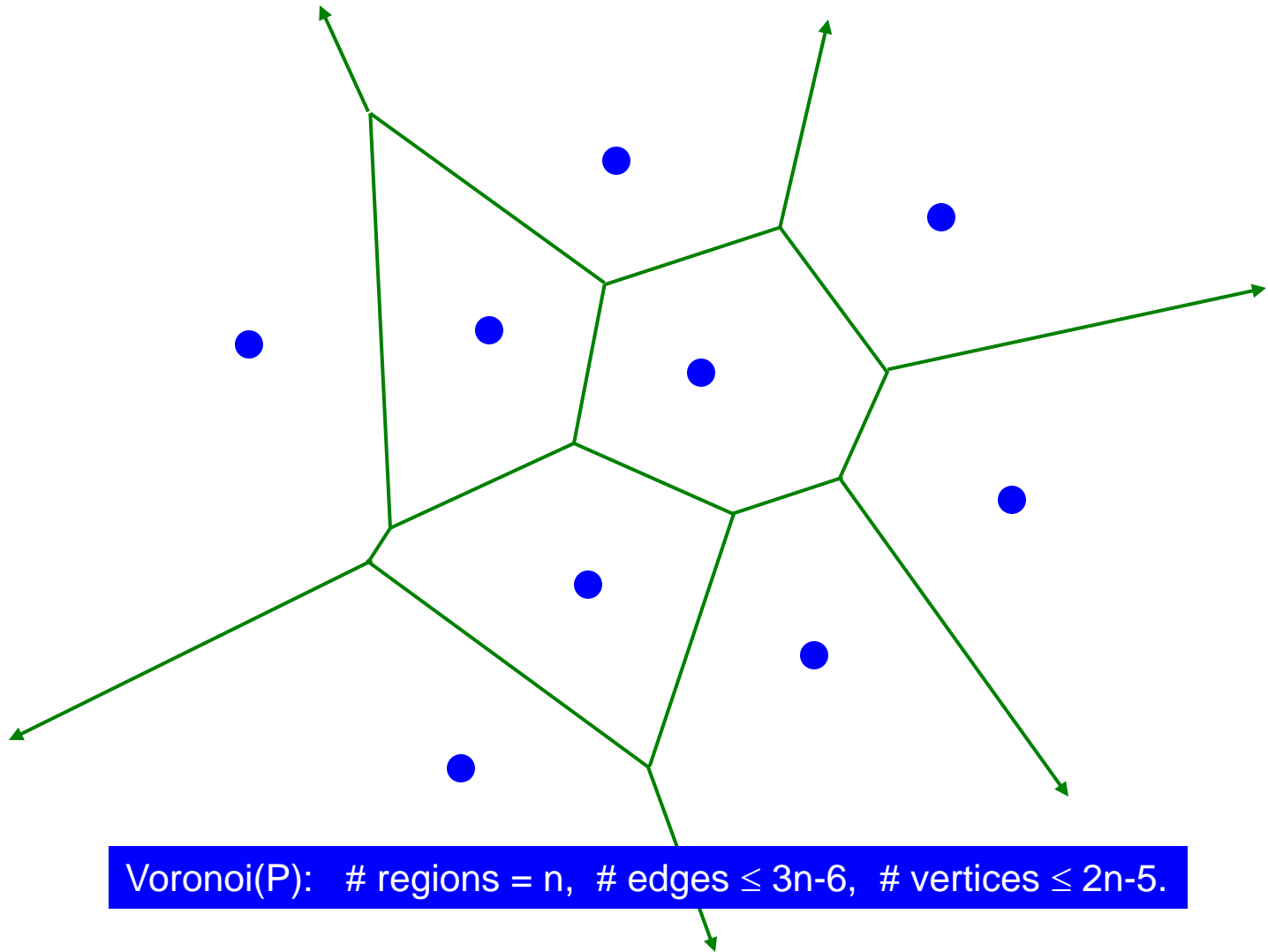
Voronoi Diagram

$P = \{ p_1, p_2, \dots, p_n \}$ a set of n points in the plane.



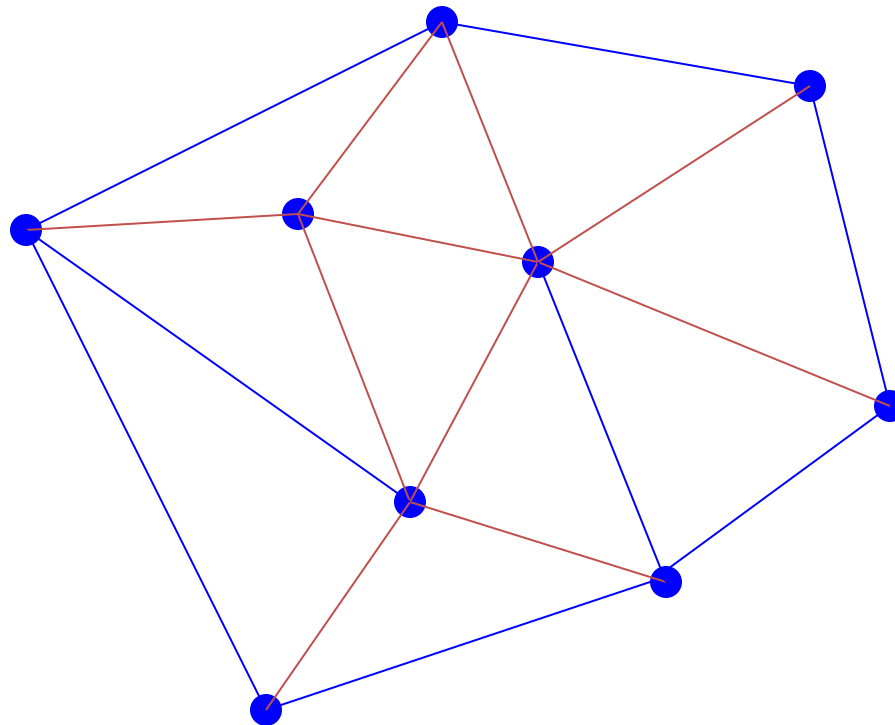
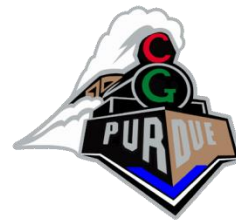


Voronoi Diagram:



Voronoi(P): # regions = n, # edges $\leq 3n-6$, # vertices $\leq 2n-5$.

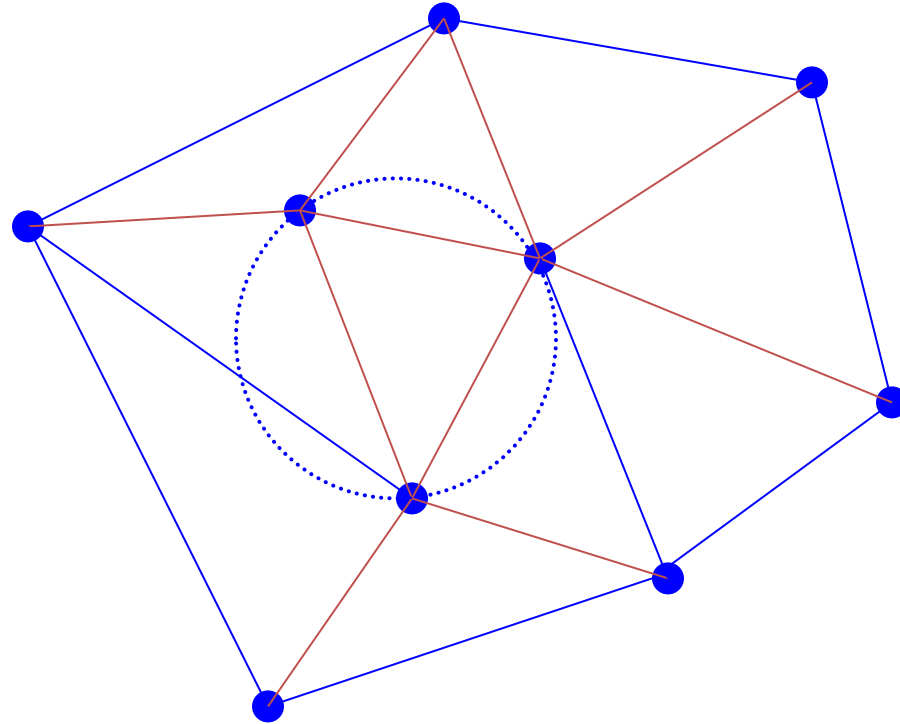
Delaunay Triangulation = Dual of the Voronoi Diagram



DT(P): # vertices = n , # edges $\leq 3n-6$, # triangles $\leq 2n-5$.



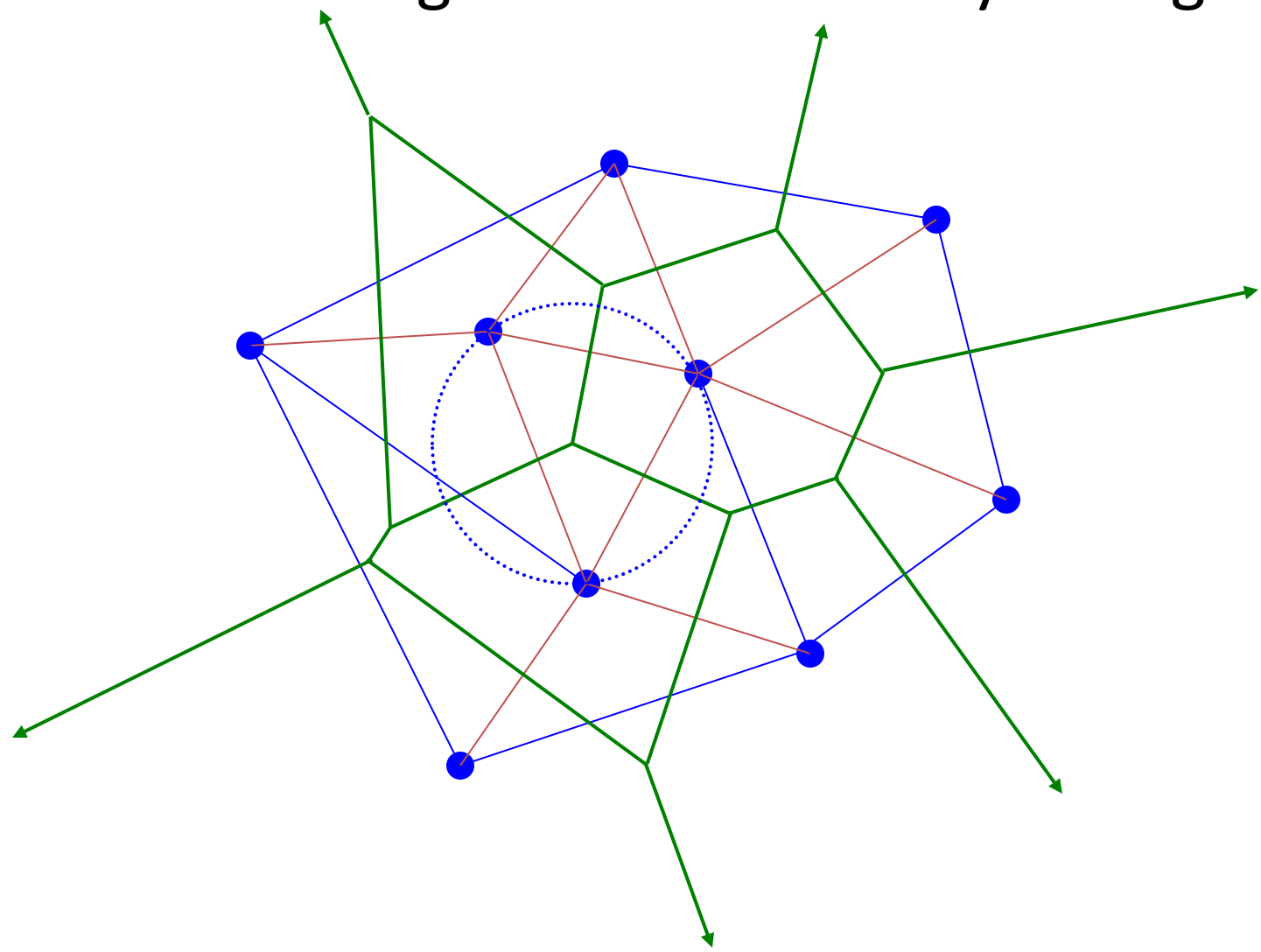
Delaunay Triangulation



Delaunay triangles have the “empty circle” property.



Voronoi Diagram and Delaunay Triangulation

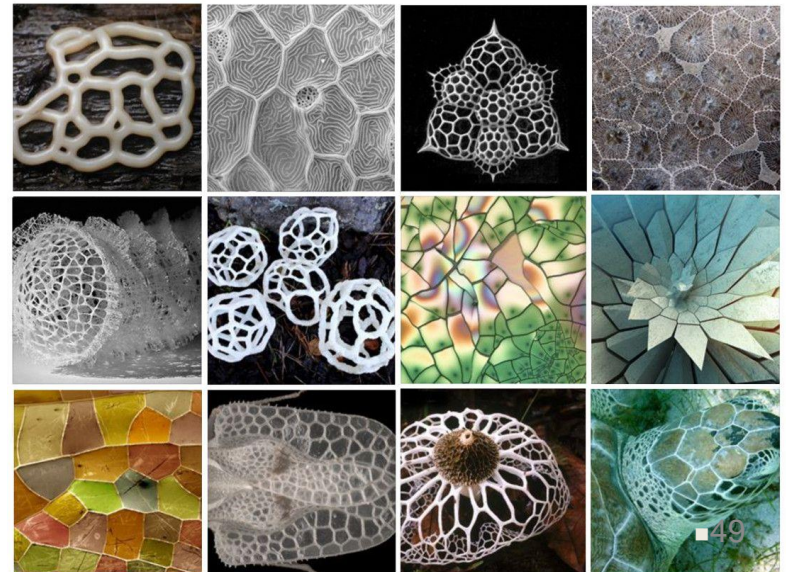




VD Properties

- Each Voronoi region $V(p_i)$ is a convex polygon (possibly unbounded).
- $V(p_i)$ is unbounded $\Leftrightarrow p_i$ is on the boundary of $\text{CH}(P)$.
- Consider a Voronoi vertex $v = V(p_i) \cap V(p_j) \cap V(p_k)$.
Let $C(v)$ = the circle centered at v passing through p_i, p_j, p_k .
- $C(v)$ is circumcircle of Delaunay Triangle (p_i, p_j, p_k) .
- $C(v)$ is an empty circle, i.e., its interior contains no other sites of P .

Voronoi Regions in Nature



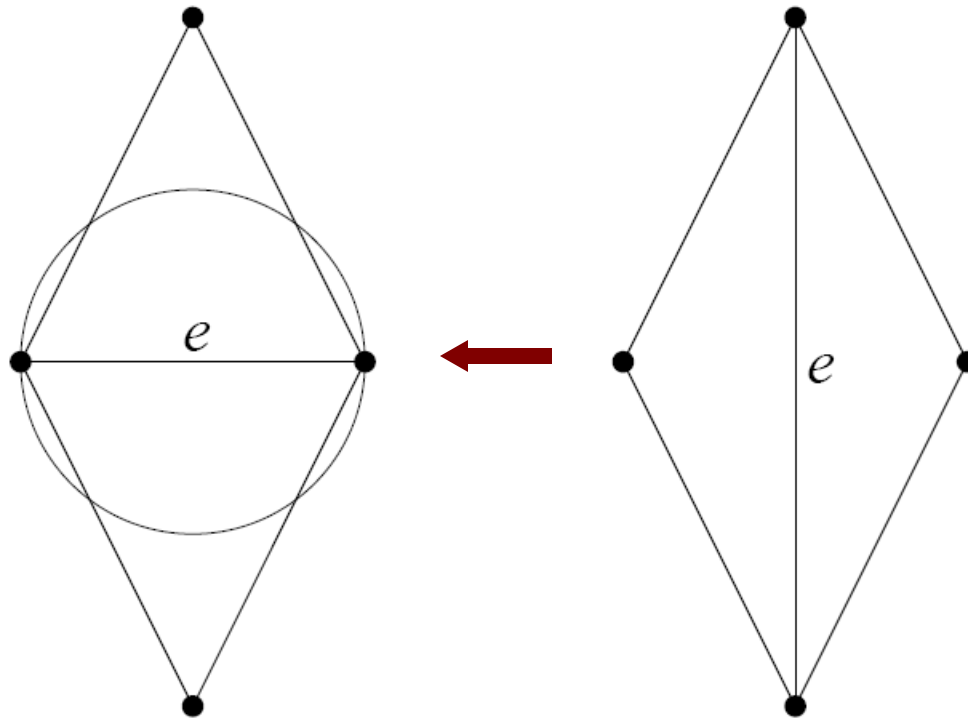


Computing Delaunay Triangulation

- Many algorithms: $O(n \log n)$
- Lets use flipping:
 - Recall: A *Delaunay Triangulation* is a set of triangles T in which each edge of T possesses at least one empty circumcircle.
 - Empty: A circumcircle is said to be empty if it contains no nodes of the set V



What is a flip?



A non-Delaunay edge flipped



Flip Algorithm

- ??



Flip Algorithm

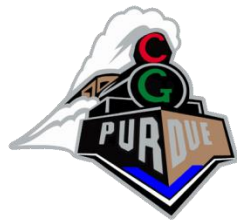
1. Let V be the set of input vertices.
2. $T =$ Any Triangulation of V .
3. Repeat until all edges of T are Delaunay edges.
 - a. Find a non-delaunay edge that is flippable
 - b. Flip

Naïve Complexity: $O(n^2)$

Locally Delaunay \rightarrow Globally Delaunay



- If T is a triangulation with all its edges locally Delaunay, then T is the Delaunay triangulation.
- Proof by contradiction:
 - Let all edges of T be locally Delaunay but an edge of T is not Delaunay, so flip it...



Flipping

- Other flipping ideas?



Randomized Incremental Flipping

- Complexity can be $O(n \log n)$



Fortune's Algorithm

- “A sweepline algorithm for Voronoi “Algorithms”, 1987, $O(n \log n)$

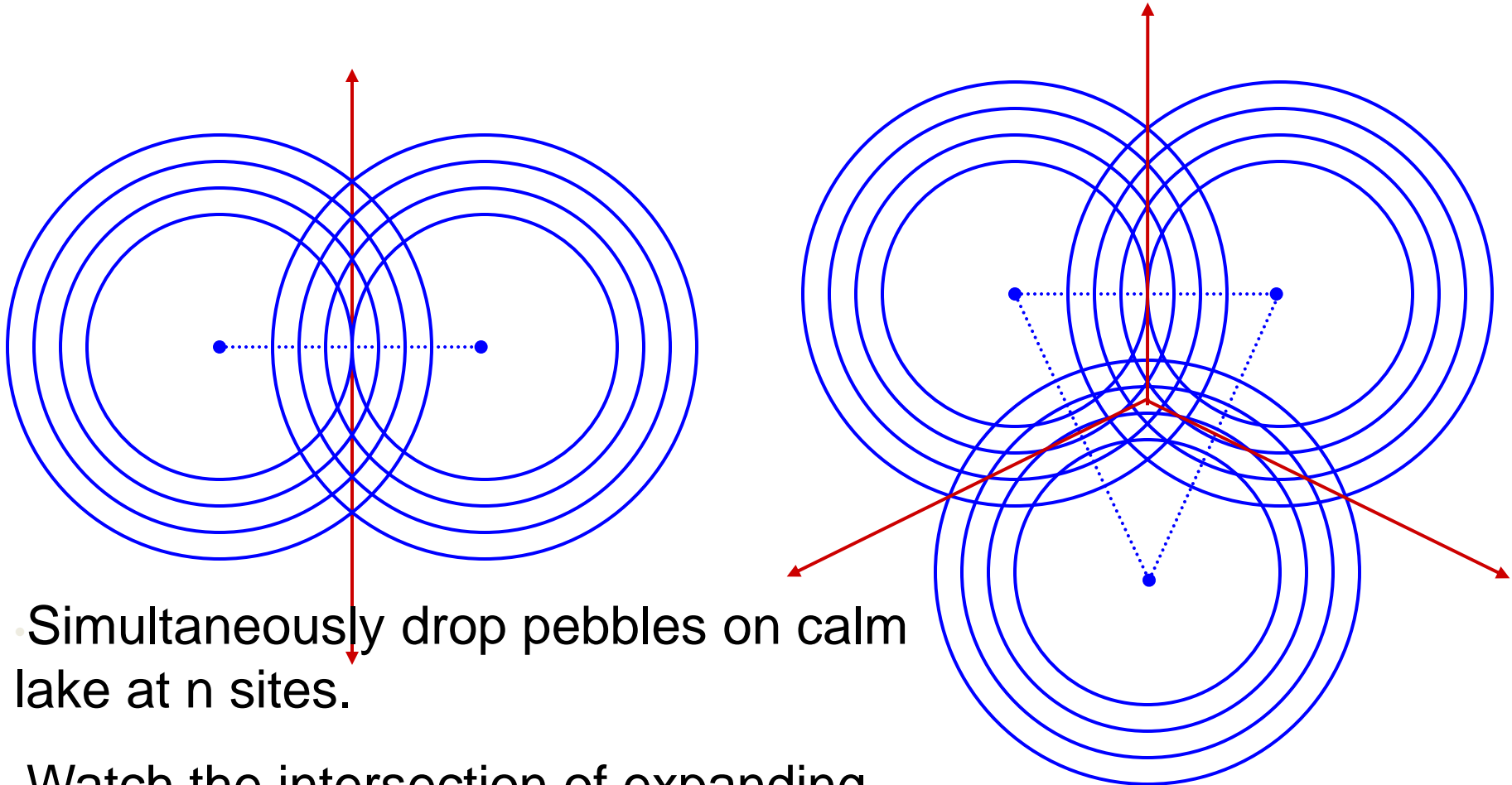
<https://www.youtube.com/watch?v=k2P9yWSMaXE>

Pseudocode:

```
add a site event in the event queue for each site
while the event queue is not empty
    pop the top event
    if the event is a site event
        insert a new arc in the beachline
        check for new circle events
    else
        create a vertex in the diagram
        remove the shrunk arc from the beachline
        delete invalidated events
        check for new circle events
```

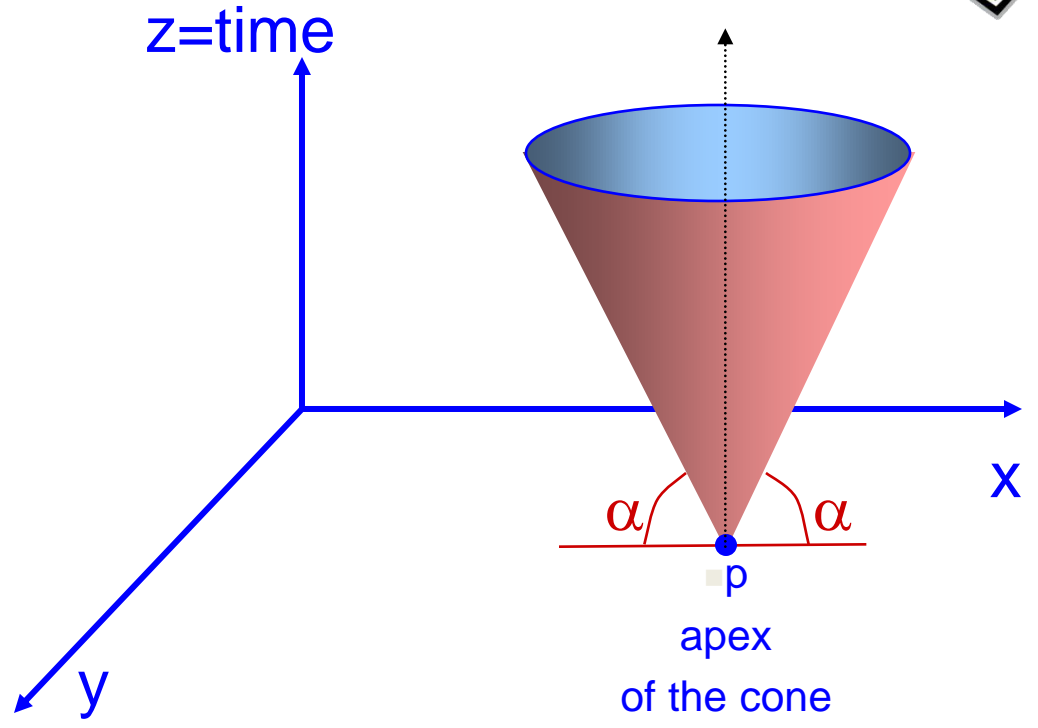
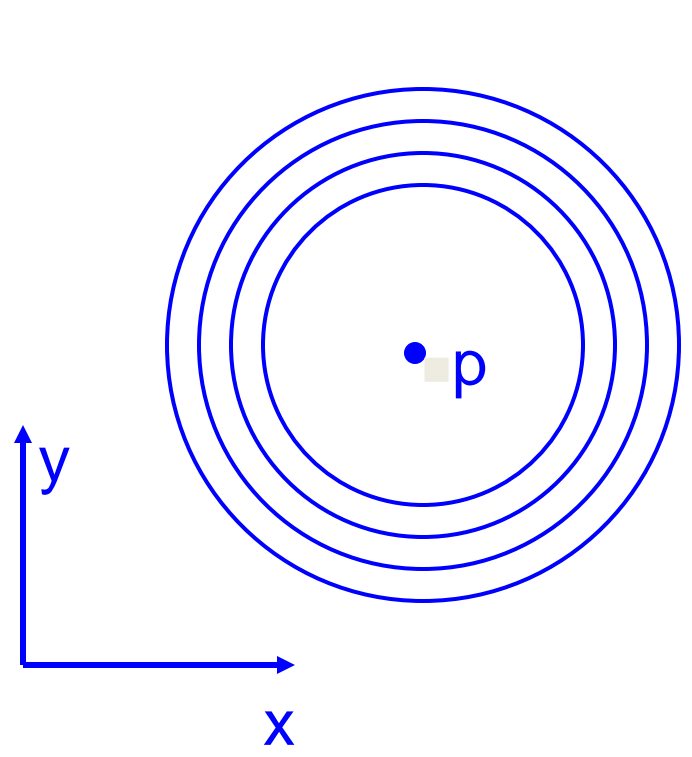


Wave Propagation View



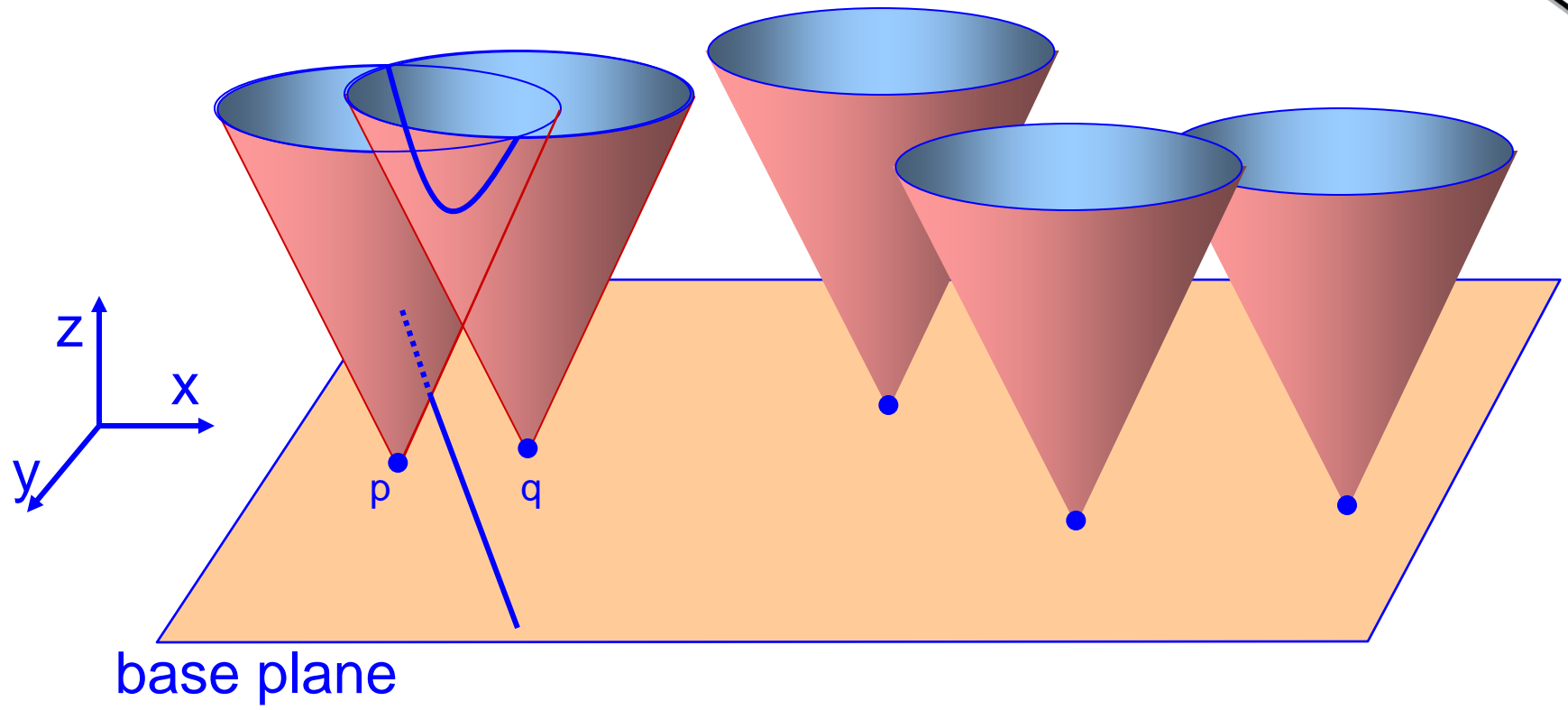
- Simultaneously drop pebbles on calm lake at n sites.
- Watch the intersection of expanding waves.

Let Time be the 3rd Dimension



All sites have identical opaque cones.

Let Time be the 3rd Dimension



All sites have identical opaque cones.

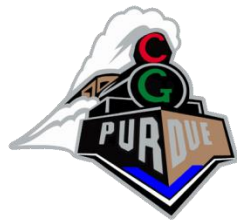
$\text{cone}(p) \cap \text{cone}(q) = \text{vertical hyperbola } h(p,q).$

vertical projection of $h(p,q)$ on the xy base plane is $PB(p,q).$



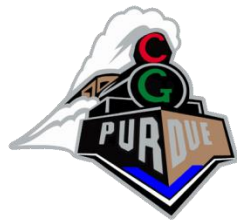
Voronoi Diagrams

- <http://alexbeutel.com/webgl/voronoi.html>

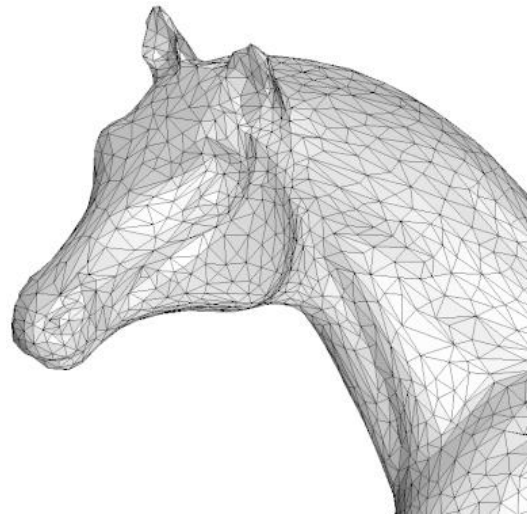
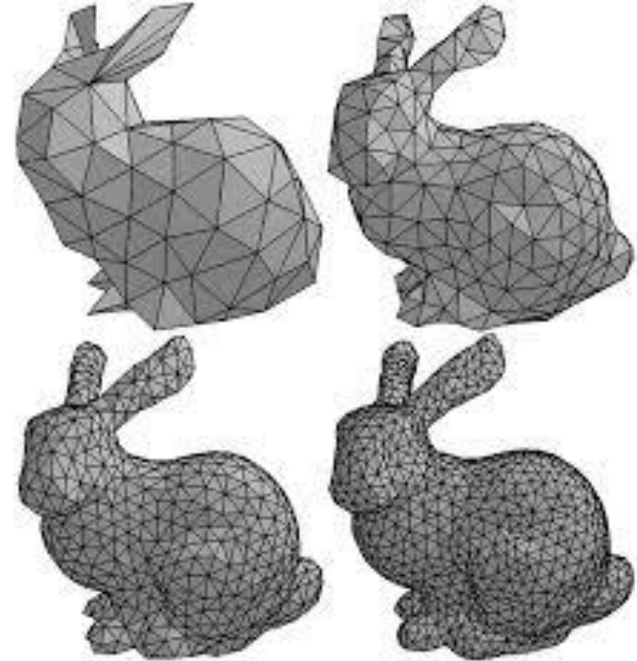
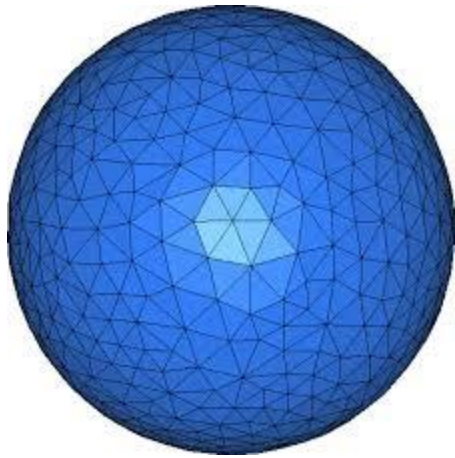


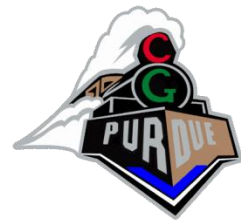
Voronoi Diagram

- <http://www.raymondhill.net/voronoi/rhill-voronoi-demo5.html>



Examples Triangulations





And Beyond...

- Not “relaxation” but more general:
 - Reaction Diffusion...
 - <https://pmneila.github.io/jsexp/grayscale/>
 - Textures:

