



CS18000: Problem Solving And Object-Oriented Programming

Types as Data Abstraction

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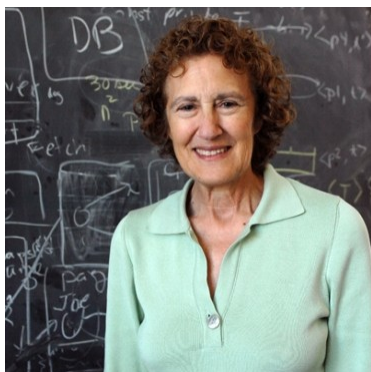
Abstraction: How Programming Scales



- Functional Abstraction
 - Clear specification of what a function does
 - Don't sweat the details of *How* something is done
- Data Abstraction
 - Clear specification of what something *is*
 - Don't worry about how it is represented



Know who this is?



- Barbara Liskov
 - 2008 ACM Turing Award winner
 - “For contributions to practical and theoretical foundations of programming language and system design, especially related to **data abstraction**, fault tolerance, and distributed computing.”

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Data Abstraction: Types



- `char ch = 'a'`
- How is `ch` *really* stored in the computer?
 1. `a`
 2. `97`
 3. `01100001`
 4. None of the above
- The `char` data type is a data abstraction

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What makes a data type?



- Semantics: *What* is being represented
 - not how
- What we can do with it
 - Functional abstractions
- Constraints
 - Can all operations be applied to any object of the type?
 - Can an object be modified?
- Special values
 - constants

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Primitive Types as Data Abstractions



Primitive Type

- Integers
 - byte, short, int, long
- Rationals
 - float, double
- Boolean
 - boolean
- Character
 - char

Class

- Number
 - Byte, Short, Integer, Long
 - BigInteger
- Number
 - Float, Double
- Boolean
- Character

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Primitive Type vs. Class



- Primitive types are “special”
 - literals (constant values)
 - infix operators (+, -)
 - ...
- Classes used to generate objects
 - double 3. vs. object of type Double containing the value 3.
- *For now, big issue is equality test*
 - `a.equals(b)` vs. `a == b`
 - `==` : variables `a` and `b` refer to the same object
 - *Two different objects can have same value!*

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Type: (semi)Formal Definition



- Domain: Set of possible values
 - Infinite?
- Operations
 - Constructor
 - Field Access
 - Methods – Accessor, Mutator
- Invariants
 - Must be true at completion of every operation

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Type: Implementation as Class



- Representation
 - Must represent all possible domain values
- Operations
 - Define Constructor(s)
 - Define methods
- Invariants
 - Ensure invariants hold after execution of Constructor(s), methods

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Field Access vs. Accessor Methods



- `Point p = new point(3.0,4.0);`
- We want the horizontal axis:
 - `x = p.x;`
 - `x = p.getx();`
- Which is better? Why?
- What if we wanted:
 - `r = p.radial;`
 - `theta = p.angle;`

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Class example: Point



```
public class Point {
    private double x; \\ Horizontal
    axis
    private double y; \\ Vertical axis
    public Point ( double px; double
    py ) {
        x = px;
        y = py;
    }
    public double getX() {
        return x;
    }
    public void setX(double x) {
        this.x = x;
    }
}

public double getR() {
    return Math.sqrt(x*x+y*y);
}
public void setTheta(Angle
theta) {
    double r = getR();
    x = r*theta.cos();
    y = r*theta.sin();
}

public Point (double r, Angle
theta) {
    this(r *theta.cos(), r*theta.sin());
}
...
}
```

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Class example: Point



```
public class Point {
    private double x; \\ Horizontal
    axis
    private double y; \\ Vertical axis
    public Point ( double px; double
    py ) {
        x = px;
        y = py;
    }
    public double getX() {
        return x;
    }
    public void setX(double x) {
        this.x = x;
    }
}

public double getR() {
    return Math.sqrt(x*x+y*y);
}
public void setTheta(Angle
theta) {
    double r = getR();
    x = r*theta.cos();
    y = r*theta.sin();
}

public Point (double r, Angle
theta) {
    this(r *theta.cos(), r*theta.sin());
}
...
}
```

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(name) Overloading



- Two methods named “Point”
 - Which is which?
- Signature: Method identified by
 - type (class),
 - name, and
 - *type/number of arguments*
- Can do for any method
 - Use sparingly, can be confusing

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Referring to Self: this



- this – object of class method is in
 - Specifically, the instance that was invoked
 - In getX(), return x;
 - x is shorthand for this.x;
- Allows parameters, instance variables to have same name
 - Makes interface (abstraction) cleaner

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Calling methods in methods



- Methods in a class can be called from other methods
 - But be careful
 - Don't call from constructor
 - Except other constructor
 - Or static method
 - *static methods can be called anytime*
 - But can't access instance variables
- Suggestion: Think Invariants
 - Method (may) assume valid instance when called
 - If you are modifying instance, does this hold?
 - If not, may not be safe

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Classes: Some Confusing Keywords



Access

- public
 - Part of “abstraction”
- private
 - Only accessible inside the class
- (default)
 - The same *package*
- protected
 - package as well as inheriting classes (we'll cover this later)

Modifiers

- static
 - Method: can be invoked without a valid instance
 - Variable: Class variable
 - Same variable shared among class instances
- final
 - Must be initialized
 - *Typically doesn't change*
- transient
 - Not saved if instance “stored”
- volatile
 - Used with threads to enforce consistency

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Goal: Separate Abstraction and Implementation



- Using an Abstract Data Type
 - Know *what* it is / *what* it does
 - *How* isn't important
- Once defined, can use
 - Doesn't need to be implemented
 - *(okay, we can't run the program – but we can write a program using it)*

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Abstract Data Type: (semi)Formal Definition



- Domain: Set of possible values
- Operations
 - Constructor
 - Field Access
 - Methods – Accessor, Mutator
- Invariants
 - Must be true at completion of every operation

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Abstract Data Type: Definition



- Challenge: How do we know we are using the type correctly?
 - May not even be written yet
 - Can we compile our program?
 - Or do we get obscure run-time errors?
- Solution: Define “unimplemented” class
 - **interface** for class

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Java interface



- Defines an (abstract) data type
 - Javaspeak: *reference* type
- Declares
 - Constants (static final variables)
 - Method type, name, and parameters
 - Assumed to be public
- No constructor(s)
 - Cannot construct an instance of an interface

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Interface Example



```
interface Queue {  
    // Constants  
    Student EMPTYRESULT =  
        null;  
    // Methods  
    void enqueue ( Person s );  
    Person dequeue();  
    boolean isEmpty();  
}
```

- Constants implicitly public static final
 - Could say “public static final Person EMPTYRESULT”, but considered poor style
- Methods implicitly public abstract
 - Again, poor style to say it
 - *Cannot* be static
 - Why?

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Using an interface



- Can define variables of type
 - Queue q;
 - Cannot create an instance, though
 - *In the above, q has the value null*

- Use as parameters

```
void bankTeller(Queue q) {  
    while ( ! q.isEmpty() ) {  
        manageCustomer(q.dequeue())  
    }  
}
```

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Implementing an interface



- Implement *all* methods defined in the interface
 - May use constants, methods, and even the interface itself
- Gives a normal class
 - Can create instances
- Instance of implementation can be given where interface used
 - variables, parameters

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Implementation Example



```
class StupidQueue implements
Queue {
    // Instance variables
    Person q[ ];
    int head, tail;
    // Constructor
    public StupidQueue() {
        head = 0;
        tail = 0;
        q = new Person[50];
        q[head] = EMPTYRESULT;
    }
    // Methods
    public void enqueue ( Person item ) {
        q[tail] = item;
        tail += 1;
        q[tail] = EMPTYRESULT;
    }
}
```

- Define representation of data
 - instance variables
- Declare and implement all methods in interface
 - Must be public

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Implementation Example: Complete



```

class StupidQueue implements
Queue {
    // Instance variables
    Person q[ ];
    int head = 0, tail = 0;
    // Constructor
    public StupidQueue() {
        q = new Person[50];
        q[head] = EMPTYRESULT;
    }
    // Methods
    public void enqueue ( Person item
    ) {
        q[tail] = item;
        q[++tail] = EMPTYRESULT;
    }

    Person dequeue() {
        return q[head++];
    }
    boolean isEmpty() {
        return q[head] ==
        EMPTYRESULT;
    }
}

```

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Implementation Example: Additional Methods



```

class StupidQueue implements
Queue {
    // Additional constants
    static final int MAXSIZE = 50;
    // Instance variables
    String q[] = new Person[MAXSIZE];
    int head = 0, tail = 0;
    // Constructor
    public StupidQueue() {
        q[head] = EMPTYRESULT;
    }
    // Methods
    public void enqueue ( Person item ) {
        if (! isFull()) {
            q[tail] = item;
            q[++tail] = EMPTYRESULT;
        }
    }

    public Person dequeue() {
        if ( isEmpty() )
            return EMPTYRESULT;
        else
            return q[head++];
    }
    public boolean isEmpty() {
        return head >= tail;
    }
    public boolean isFull() {
        return tail >= MAXSIZE;
    }
}

```

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