

Topic 8: Fast Fourier
Transform

Polynomial Multiplication

$$f(x) = 1+x \quad g(x) = 1+x+x^2 \quad f \cdot g(x) = 1+2x+2x^2+x^3$$

Naive $O(n^2)$

Observation: A poly of n coefficients is determined by n points

$$f(1) = 2 \quad f(2) = 3 \quad f(3) = 4 \quad f(4) = 5$$

$$g(1) = 3 \quad g(2) = 7 \quad g(3) = 13 \quad g(4) = 21$$

$$f \cdot g(1) = 6 \quad f \cdot g(2) = 21 \quad f \cdot g(3) = 42 \quad f \cdot g(4) = 105$$

$$\downarrow$$
$$f \cdot g(x) = 1+2x+2x^2+x^3$$

Problem: compute $f(x)$ is $O(n) \Rightarrow O(n^2)$ for n points

777

assume $n \geq k$

Intuition: We are going to find n points fast in $O(n \log n)$

$F = (f(p^0), f(p^2), \dots, f(p^{n-1}))$ for some p

$$\text{Let } f(x) = a_0 + a_1 x + a_2 x^2 + \dots + a_{n-1} x^{n-1}$$

$$= \underbrace{(a_0 + a_2 x^2 + \dots + a_{n-2} x^{n-2})}_{\text{subproblems of size } \frac{n}{2}} + x \underbrace{(a_1 + a_3 x^2 + \dots + a_{n-1} x^{n-2})}_{\text{subproblems of size } \frac{n}{2}}$$

subproblems of size $\frac{n}{2}$

$$= f_e(x^2) + x f_r(x^2)$$

$$F_e = (f_e(p^0), f_e(p^2), \dots, f_e(p^{n-2}))$$

$$F_r = (f_r(p^0), f_r(p^2), \dots, f_r(p^{n-2}))$$

$$F_i = f(p^i) = f_e(p^{2i}) + p^i \cdot f_r(p^{2i}) = F_{e,i} + p^i \cdot F_{r,i} \quad (0 \leq i < \frac{n}{2})$$

What should we do when $i \geq \frac{n}{2}$?

'Magic Numbers'

Goal: Find magic numbers p . s.t. $p^n = 1$.

Then $f_c(p^{n+i}) = f_c(p^i)$ and problem solved

$$FFT: e^{\pi i} = -1 \Rightarrow p = e^{\frac{2\pi i}{n}} = \cos\left(\frac{2\pi}{n}\right) + \sin\left(\frac{2\pi}{n}\right)i.$$

(Precision issues in CP...)

NTT:

(Fermat's little Thm) $p^{u-1} = 1 \pmod{u}$ for any p & u prime

(1) Make $u-1 = \alpha \cdot n$. then $(p^\alpha)^n = 1$

$$SP8244353 = \underbrace{2^{23}}_n \cdot 7 \cdot 17 + 1$$

(2) Make $p^0, p^1, p^2 \dots p^{u-2}$ all different

3 sets of
prime / generator
in the
reference

Known as primal root / generator

For SP8244353,
3 is a generator

Reverse ?

(Almost) the same (see reference)

How to FFT in Practice

'Generating Function':

$$f(x) = a_0 + a_1x + a_2x^2 + \dots$$

x^i is not a value, but an abstract 'item' i

And a_i is the number of ways to obtain this 'item' i .

Lightsabers

Lightsabers

take 0 red sabers

2 red sabers

$$(1 + x + x^2)$$

take 2 red sabers

1 blue saber

$$(1 + x)$$

$$1 + 2x + \boxed{2x^2} + x^3$$

2 ways to take 2 sabers

mod 100P?

observe that 1 poly mult will be at most

$$n \cdot 100P < PPF > 44353$$

so use PPF > 44353 gives accurate result

and we can clamp down by mod 100P.

Nikita and Order Statistics

Nikita and Order Statistics

s_i : number of elements in $a[0..i]$ that is $< x$

e.g. $a = (1, 2, 3, 4, 5)$ $x = 3$
 $s = (0, 1, 2, 2, 2)$

$[l..r]$ has k elements $< x \iff s_r - s_l = k$

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$$s = (\underline{0}, \underline{1}, \underline{2}, \underline{2}, \underline{2}, \underline{2}) \Rightarrow f(x) = 1 + x + 4x^2 \quad \leftarrow \text{pick } r$$
$$g(x) = 1 + x^{-1} + 4x^{-2} \quad \leftarrow \text{pick } l$$

$$f \cdot g(x) = \dots + 18 + 5x + 4x^2$$
$$(x^2 \cdot g(x) = x^2 + x + 4)$$

Largest = $n^2 > 999 \approx 44333$?

Use two primes and Chinese Remainder Theorem
or 7-7-7 and pray for precision.