

## INTRODUCTION

### What is a computer network?

Components of a computer network:

- hosts (PCs, laptops, handhelds)
- routers & switches (IP router, Ethernet switch)
- links (wired, wireless)
- protocols (IP, TCP, CSMA/CD, CSMA/CA)
- applications (network services)
- humans and service agents

Hosts, routers & links form the *hardware* side.

Protocols & applications form the *software* side.

Protocols can be viewed as the “glue” that binds everything else together.

A physical network:



Protocol example: low to high

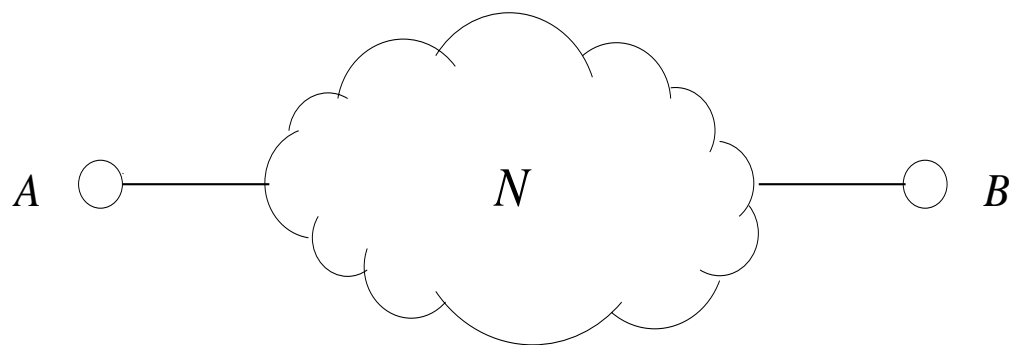
- NIC (network interface card): hardware
  - e.g., Ethernet card, WLAN card
- device driver: part of OS
- ARP, RARP: OS
- IP: OS
- TCP, UDP: OS
- OSPF, BGP, HTTP: application
- web browser, ssh: application
  - multi-layered glue

What is the role of protocols?

→ facilitate communication or networking

Simplest instance of networking problem:

Given two hosts  $A$ ,  $B$  interconnected by some network  $N$ , facilitate communication of information between  $A$  &  $B$ .



Information abstraction

- representation as objects (e.g., files)
- bytes & bits
  - digital form
- signals over physical media (e.g., electromagnetic waves)
  - analog form

Minimal functionality required of  $A$ ,  $B$

- encoding of information
  - decoding of information
- data representation & a form of translation

Additional functionalities may be required depending on properties of network  $N$

- information corruption
  - $10^{-9}$  for fiber optic cable
  - $10^{-3}$  or higher for wireless
- information loss: packet drop
- information delay: like toll booth, airport
- information security

Network  $N$  connecting two or more nodes can be

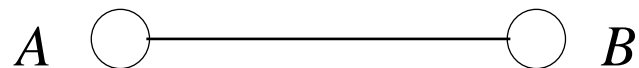
- point-to-point links
- multi-access links
- internetworks
  - physical vs. logical topology
  - e.g., peer-to-peer, VPN

Network medium may be

- wired
- wireless

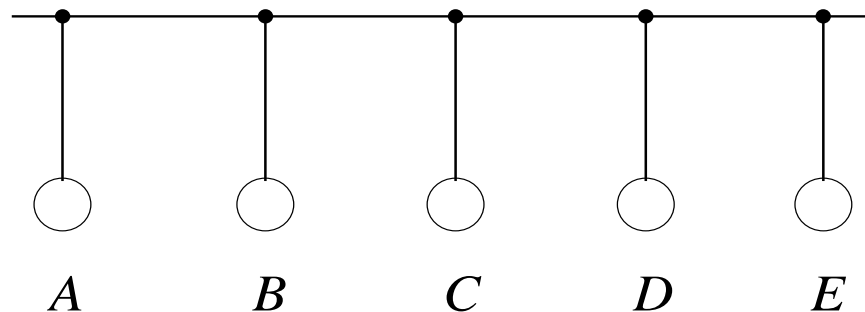
Node (e.g., hosts, routers) may be

- stationary
- mobile

*Point-to-point links*

- various “cables”
- line of sight wireless communication
  - directional antennas
- no addressing necessary
  - special case

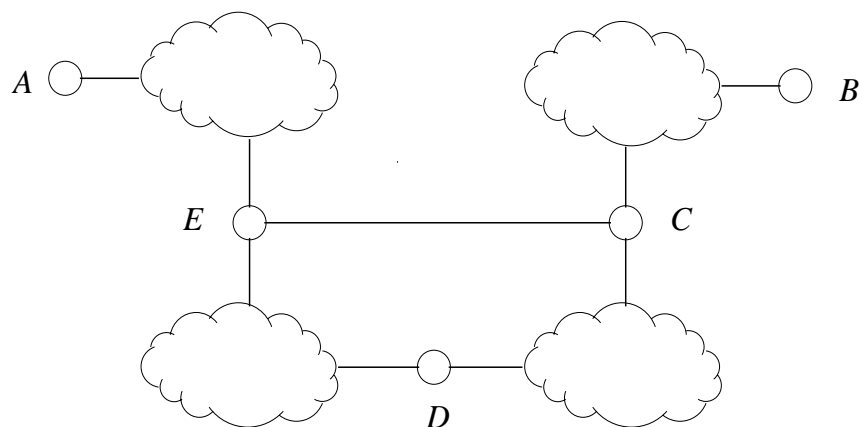
## *Multi-access links*



- bus (e.g., old Ethernet)
- wireless media
  - omni-directional antennas
- broadcast mode (physical; not logical)
  - listen to everything: promiscuous mode
- access control: i.e., bus arbitration
  - resolve contention and recover from interference
- addressing necessary



## Internetwork



- recursive definition
  - point-to-point and multi-access: internetwork
  - composition of one or more internetworks
- addressing necessary
- path selection between sender/receiver: routing
- how much to send: congestion control
- protocol translation: internetworking
- location management: e.g., Mobile IP

LAN (local area network) vs. WAN (wide area network)  
distinction:

- LAN: point-to-point, multi-access
- WAN: internetwork
  - geographical distinction is secondary
  - often go hand-in-hand
  - counter example?

Myriad of different LAN technologies co-existing in a WAN. For example:

- Fast Ethernet (100 Mbps)
- Gigabit Ethernet (1000 Mbps)
- WLAN (54 or 11 Mbps)
- FDDI (Fiber Distributed Data Interface)
- wireless Ethernet (11 Mbps, 54 Mbps)
- SONET
- ATM
- modem/DSL

→ WAN is a collection of LANs

Each LAN, in general, speaks a different language.

- message format
- procedural differences

Internetworking solves this problem by translating everything to IP ...

- technical definition of **I**nternet

But:

- IP is not necessary
- e.g., large systems of layer 2 switches
- trend: L2 (70s & 80s) → IP (90s) → L2 (Y2K+)
- IP remains central glue

Remark on addresses (aka names):

Communicating entities are *processes* residing on nodes *A* and *B* running some operating system.

Thus an *address* must also identify which process a message is destined for on a host.

→ e.g., port number abstraction

## Key Issues

### Fault-tolerance

The larger the network, the more things can go wrong.

E.g.: link/node failures, message corruption, software bugs

→ managing downtime: tier-1 providers

→ 99.999%

Two types of failures:

- independent
- correlated

In a network system with  $n$  components, assume a component fails with independent probability  $p$

- expected number of failures:  $n \cdot p$
- probability of no failures:  $(1 - p)^n$
- probability of  $k$  simultaneous failures:  $p^k$

Thus correlated failures have miniscule probability.

- exponentially small in  $k$

In reality, failures are not independent.

→ e.g., power outage, natural disasters

We have:

→ Murphy's Law

- issue of reliable communication
- reliable network services
  - main principle: redundancy
- Examples:
  - routing of messages: alternate/back-up routes
  - domain name servers: duplication
  - transmission by space probes: forward error correction (FEC)
    - also used for multimedia traffic



## Network security

Features:

- confidentiality: encryption
- integrity: message has not been tampered
- authentication: sender really is who she claims to be
  - “CIA”
  - foundation: cryptography
  - end-to-end
  - networking problem?

Modern security vulnerabilities:

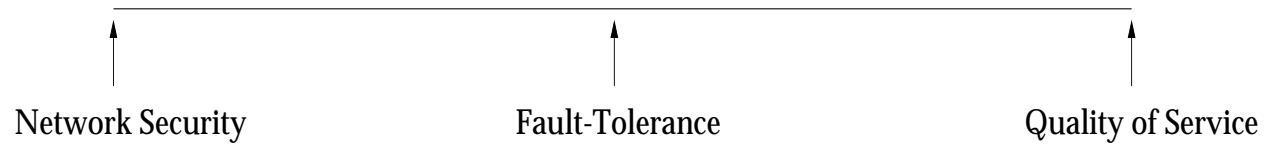
- denial of service (DoS) attack
  - e.g., SYN flooding
- distributed DoS (DDoS) attack
  - e.g., commercial, personal, infrastructure
- worm attacks: e.g., CodeRed, Blaster, ...
  - buffer overflow: mainly bugs in MS DLLs
- spam mail (security issue?)

- with fault-tolerance impacts QoS (quality of service)
  - Aug. 04: US broadband deployment exceeds dial-up
- security: trade-off with overhead
  - what is the desired operating point?
  - too much  $\Rightarrow$  too slow
  - too little  $\Rightarrow$  too vulnerable

For example: secure routing (S-BGP)

→ “BBN vs. Cisco”

Big picture:



- points in the same spectrum
- malicious (Byzantine) vs. non-malicious
- availability
- service assurances

## Performance

Issues:

- excessive traffic can cause congestion (analogous to highways)
- traffic volume exhibits large fluctuations
  - burstiness
- multimedia traffic is voluminous (even for single user)
- ubiquitous access
  - wired/wireless Internet

Potential for bottleneck development

- spontaneous or persistent
- similar consequences as failures

Different applications require different levels of service quality.

Challenges:

- how to provide customized QoS
- many users and applications: scalability
- must interoperate with legacy Internet

Current state:

- overprovisioning
  - “throw bandwidth at the problem”
  - tier-1 ISPs use sophisticated traffic engineering
- still no Internet QoS
  - changing with VoIP and content deployment
- not economic
  - few tier-1 providers make money

## Data networking, telephony, and content convergence

→ Y2K+ trend

- VoIP (Voice-over-IP): wired world

→ traditional TDM-based telephony system is entirely separate network

→ economic factors are dictating convergence

→ from KaZaA to Skype

- cellular voice networks: 2G, 2.5G, 3G

→ what is 4G?

→ telcos/cellular providers are concerned

→ take-over by WLAN + IP?

→ strategy: active participation

- peer-to-peer: rampant content dissemination
  - from audio to movies
  - content providers need to get into the action
  - do not want to get into the action

\$6 question:

→ what will the wireless/wireline future hold?

Mixture of high bandwidth/low bandwidth networks, wireline/wireless, . . .