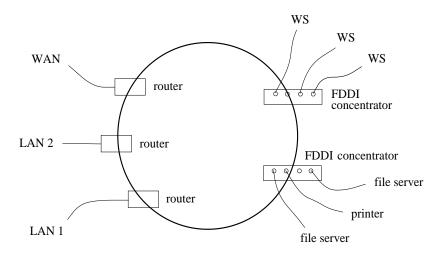
$\longrightarrow$  token ring architecture

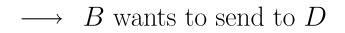
High-bandwidth extension of IBM 4 Mbps token ring and 16 Mbps IEEE 802.5 token ring standard.

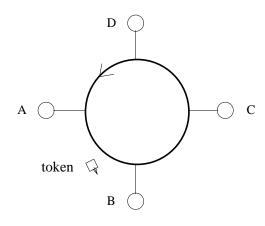
 $\rightarrow$  100 Mbps bandwidth

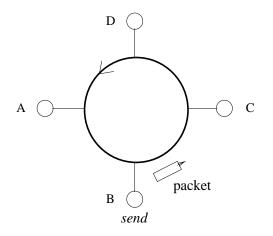
Used as high-bandwidth campus/city backbone.

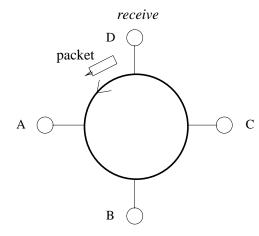
 $\rightarrow$  metropolitan/campus distance: MAN

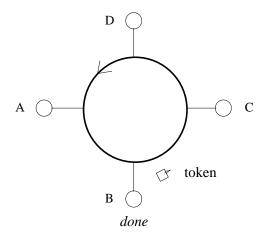




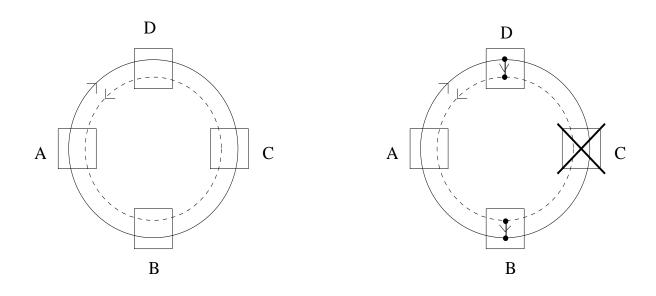




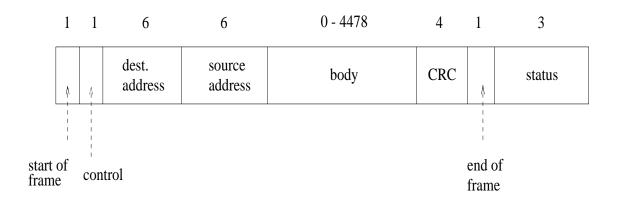




## Fault-tolerance:



- DAS (dual attachment station)
- SAS (single attachment station)



- frame size < 4500 B
- 4B/5B encoding
- synchronous/asynchronous data
- $\bullet~2~{\rm km}$  inter-station distance
- 200 km diameter (multimode fiber); 100 km circumference

Performance issues: fairness and efficiency

- TRT (token rotation time)
- THT (token holding time)

 $TRT = no. of nodes \times THT + link latency$ 

To increase efficiency: increase THT

- $\longrightarrow$  let station send as much as it needs
- $\longrightarrow$  same as frame size  $\uparrow$
- $\longrightarrow$  THT  $\uparrow \implies \rho \uparrow$

To increase fairness: limit THT

 $\longrightarrow$  limit station's one-time sending of data

To facilitate fairness: introduce TTRT (target token rotation time).

THT determining factor (assume TTRT is given):

- prioritized frames: synchronous/asynchronous
- Synchronous frames always get sent.
- If TRT > TTRT, then late; don't send asynchronous data.
- If TRT ≤ TTRT, then early; send asynchronous data for max { TTRT − TRT, single frame time } duration.

How to set TTRT?

- $\longrightarrow$  token claim process
- $\longrightarrow$  initiate when needed (e.g., start-up)
- Each station submits claim frame containing TTRT bid.
- Smaller TTRT bid overrides higher TTRT bids.
  - Compare claim frame bid against own desired TTRT.
  - If less, then reset own TTRT to lower value.
  - If larger, then put lower bid on claim frame and forward.
- Winner: same bid value when claim frame makes full circle.

 $\longrightarrow$  leader election

At the end of the day, consistent TTRT value among all stations.

 $\longrightarrow$  consensus problem

Compare against Ethernet's CSMA/CD.

- $\longrightarrow$  round-robin reservation
- $\longrightarrow$  absence of MA and collision
- $\longrightarrow$  determinism vs. indeterminism
- $\longrightarrow$  imperfect QoS assurance
- $\longrightarrow$  performance vis-à-vis CSMA/CD?

Cooperative vs. noncooperative protocols

- $\longrightarrow$  robust if some users use selfish MAC
- $\longrightarrow$  could be malicious